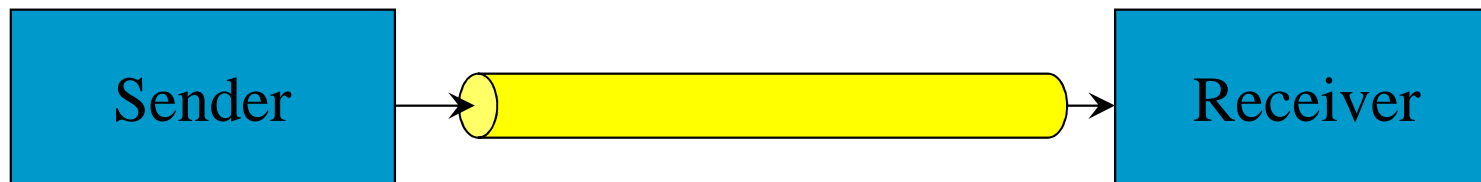


Réseaux

Channels and Multiplexing

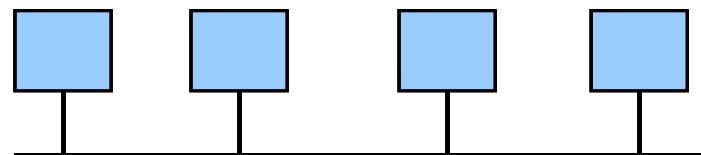
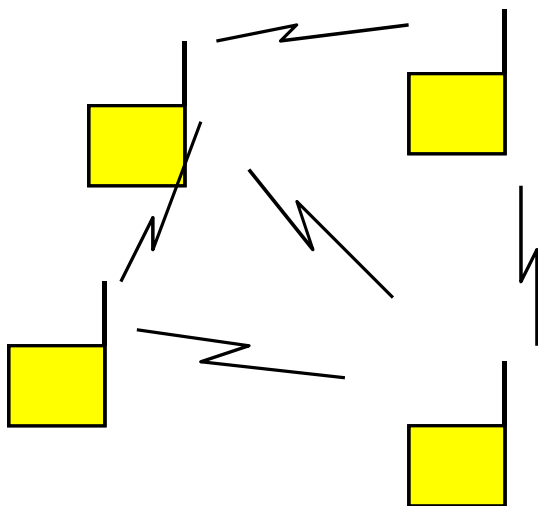
Point-to-point Channels

- They are permanent connections between a sender and a receiver
- The receiver can be designed and optimized based on the (only) signal it must receive
- Data transmission can be *continuous* or divided into *frames* (this raises synchronization problems)



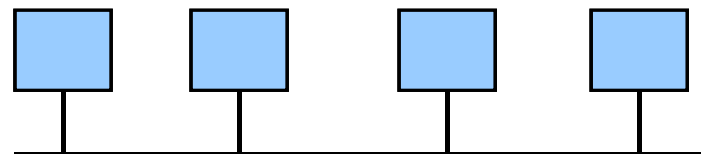
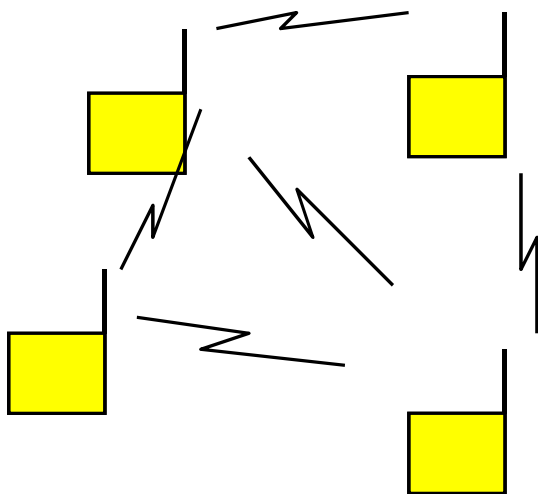
Broadcast Channels

- Many stations/nodes can access a broadcast channel in parallel. The channel is «shared» among all stations
- The transmissions of a station reaches all other stations



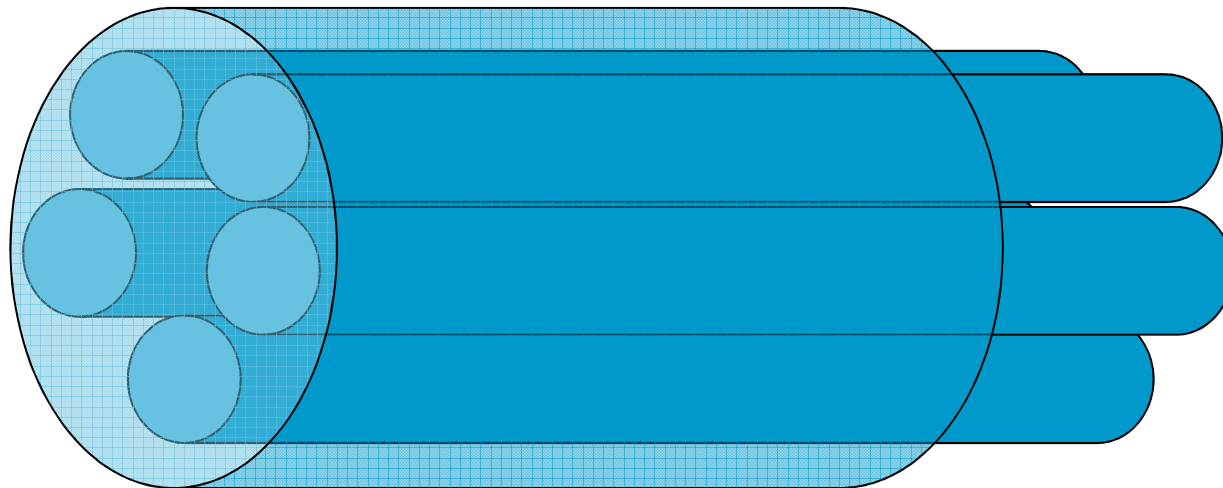
Broadcast Channels

- The receiver can receive several transmissions which differ in their *power level* and *synchronization*, and it must be able to adapt itself to such differences, and single out the right transmission
- Transmissions usually start with a *preamble* (*synchronization character*) to achieve synchronization
- Examples: local area networks/ethernet, cellular systems



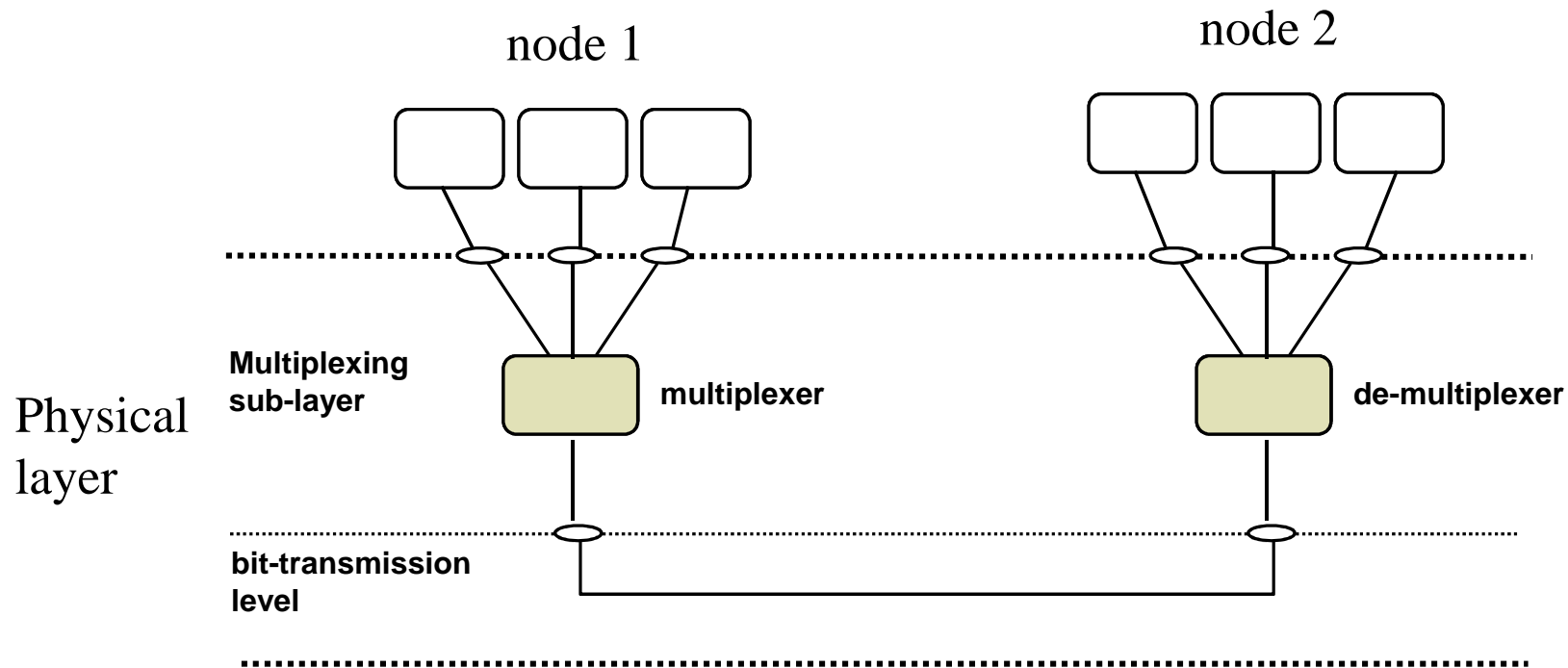
Multiplexing

- **The physical capacity of a channel can be subdivided to obtain more (sub)channels with lower speed.**



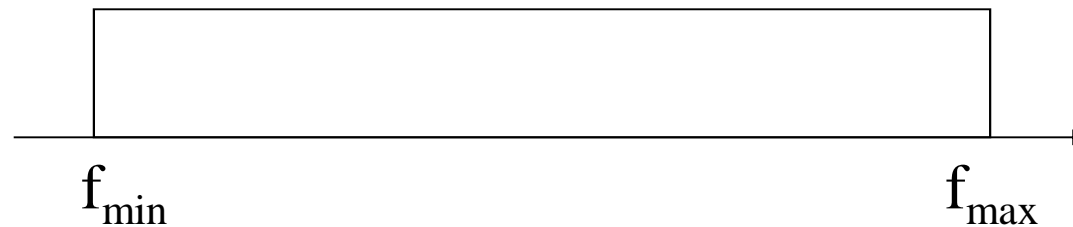
Physical Multiplexing

- Each (sub)channel is defined based exclusively on physical parameters, like frequency, time, code, wavelength ...

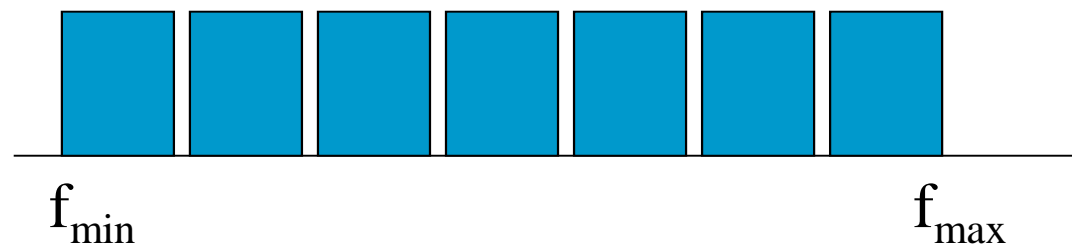


FDM (Frequency Division Multiplexing)

- Each physical channel can be characterized by its available bandwidth (the set of frequencies available for transmission, from f_{\min} to f_{\max})

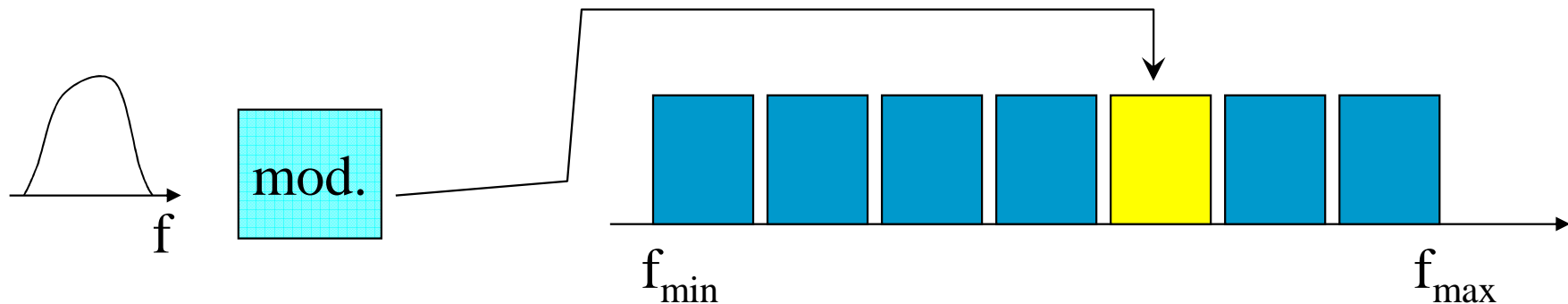


- Such bandwidth can be divided into sub-channels, and we can associate a communications to each sub-channel

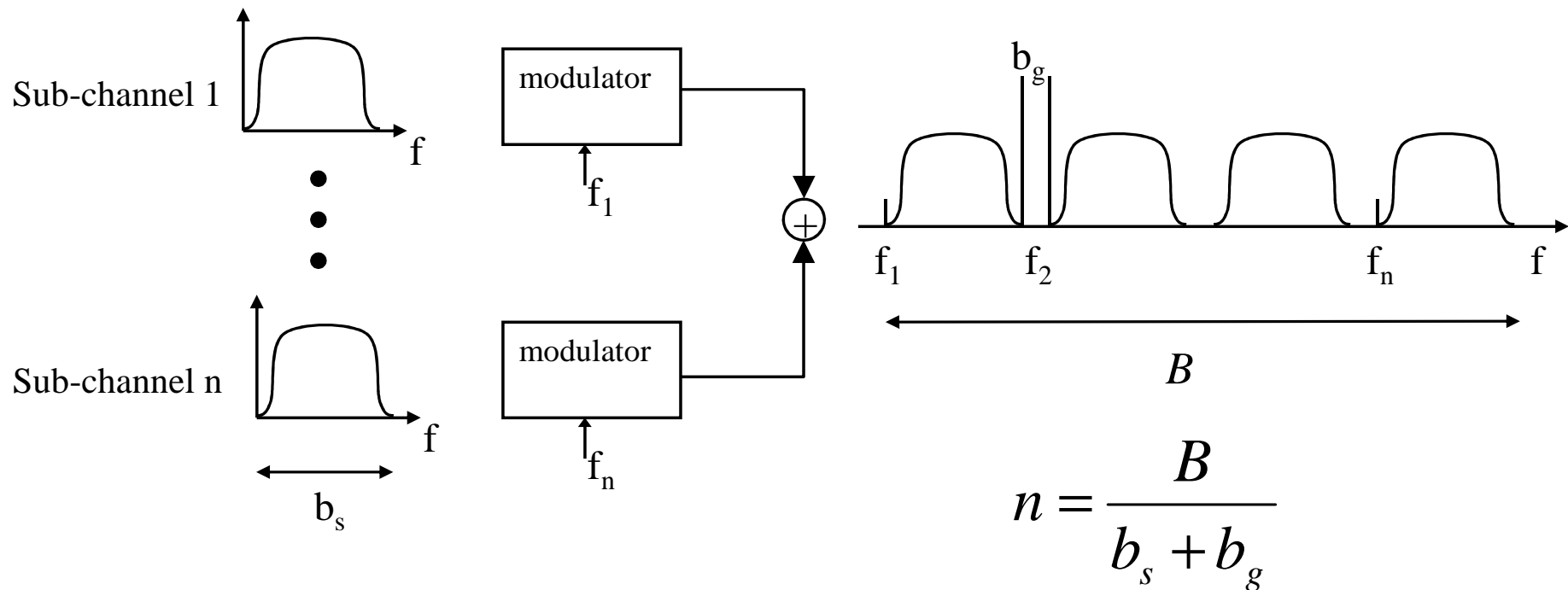


FDM (Frequency Division Multiplexing)

- The signal related to one communication is *filtered* and then *modulated* (hence, shifted in frequency) in order to fit exactly into one sub-channel



FDM (Frequency Division Multiplexing)



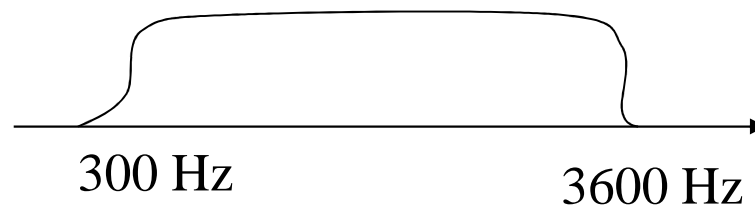
$B = \text{total available bandwidth } (f_{\max} - f_{\min})$

$b_s = \text{signal bandwidth}$

$b_g = \text{guard band}$

FDM - Telephony

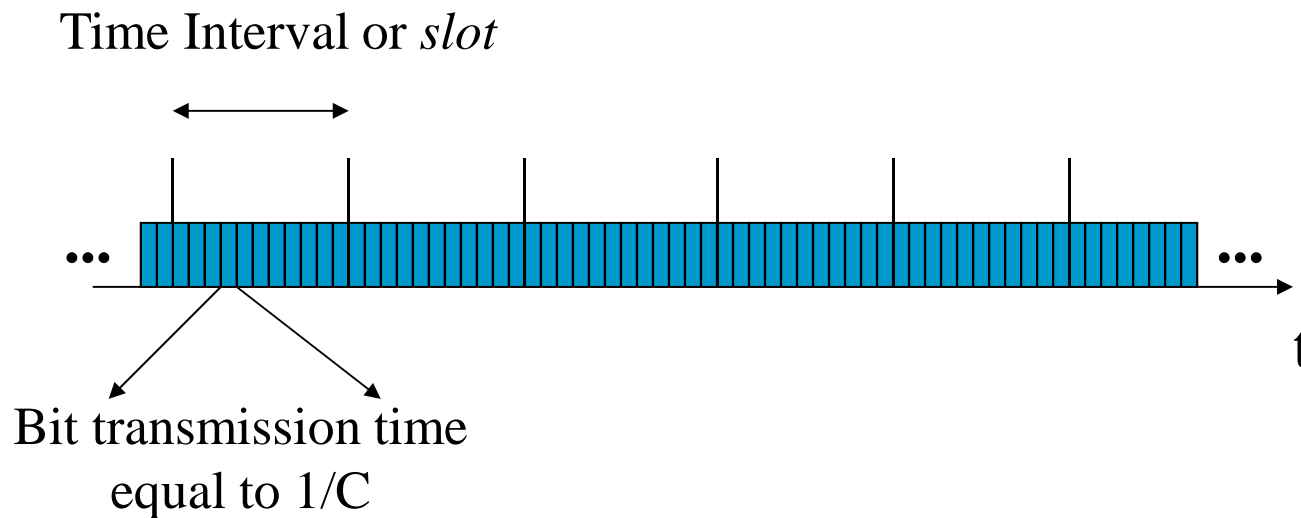
- In the past, FDM was used as a multiplexing technique to transmit voice calls between phone centrals
- Voice call bandwidth: approximately 4 kHz



- 12 channels/voice calls of 4 kHz each were multiplexed over a total bandwidth of 48 kHz (in the range between 60 and 108 kHz)
- Then, such 48 kHz aggregation was further multiplexed in even larger aggregations (in a hierarchical modulation scheme)

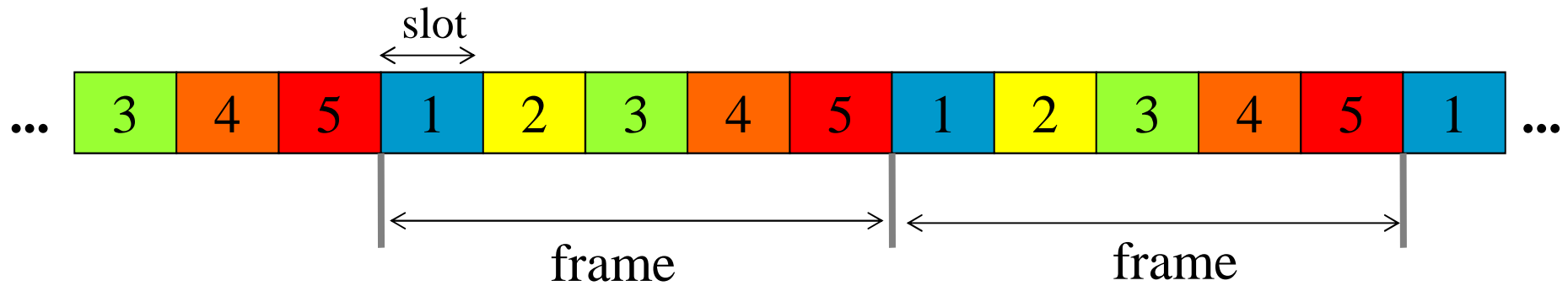
TDM (Time Division Multiplexing)

- This technique is used for digital/binary signals (sequences of 0s - 1s)
- Given a channel with speed/capacity C (bit/s), we define time intervals (named *slots*), whose duration is a multiple of the bit duration $t_b=1/C$

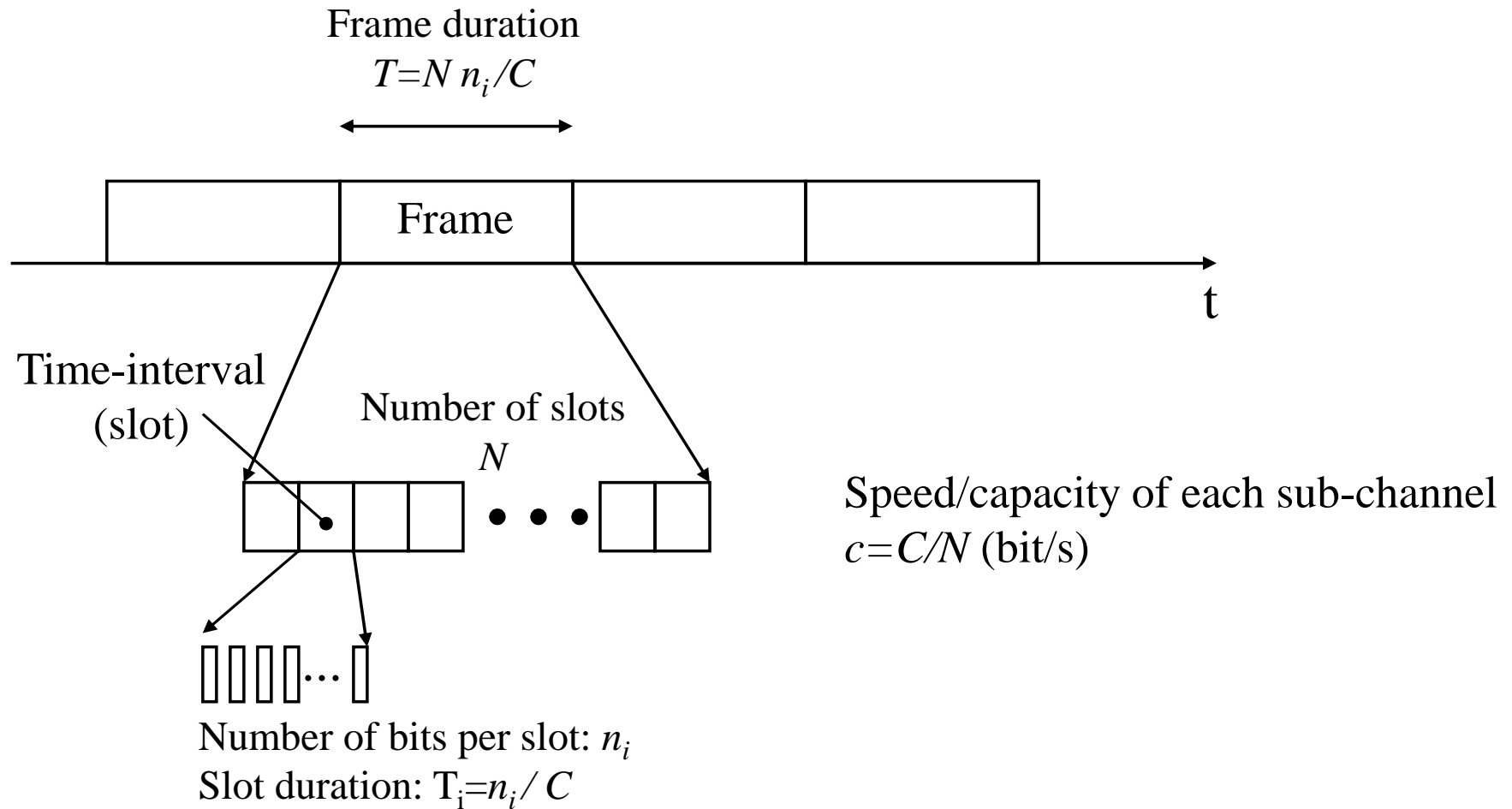


TDM (Time Division Multiplexing)

- Each source/sender can use only a single time slot every N
- Hence we define a *frame structure*, where the frame is constituted by N consecutive time slots
- If we give a number to each time slot, each source/sender is associated to a time-slot number, and it can transmit only *inside* such slot

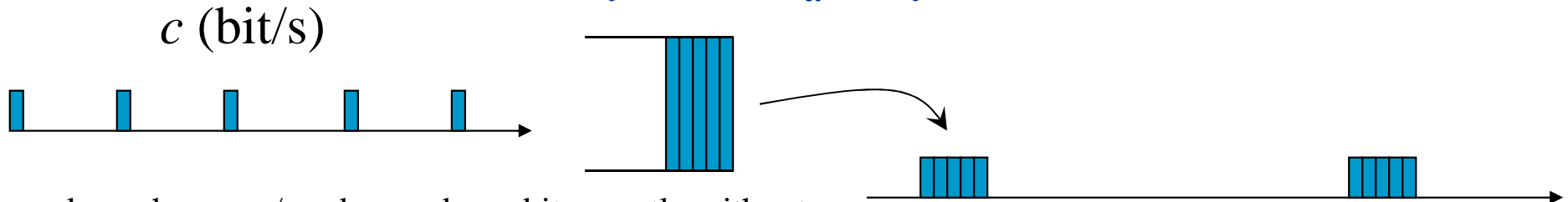


TDM (Time Division Multiplexing)



TDM (Time Division Multiplexing)

- The choice of the *slot duration* is very important (this is a parameter chosen when the slotted system is designed):
 - n_i number of bits per slot
 - T_i slot duration ($T_i = n_i / C$)
- the sub-channel capacity/speed c does not depend on T_i but only on N ($c = C/N$)
- Time to collect n_i bits: $T_a = n_i / c$



Remark: each source/sender produces bits exactly with rate c
The n_i bits that fit in the slot **must** be already available when the slot begins.
Clearly, the source needs n_i / c seconds to produce and accumulate the n_i bits

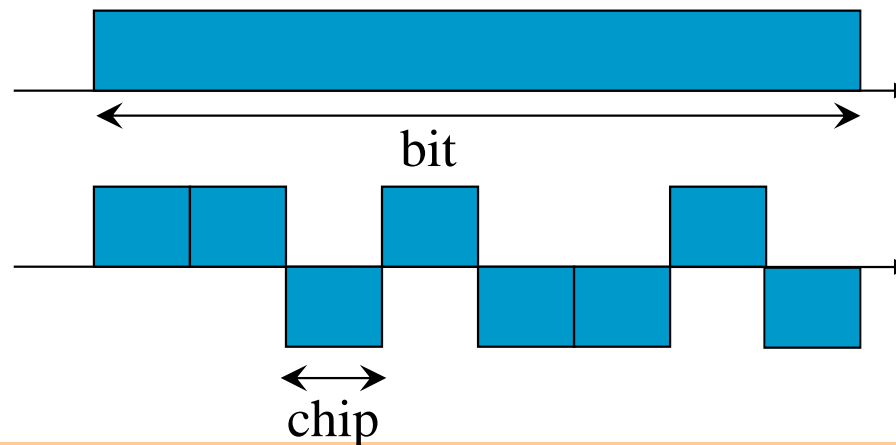
Exercise

- Let us consider a channel with capacity $C=900$ kbit/s
- We want to create 5 sub-channels: 4 with capacity $c=200$ kbit/s and 1 with capacity 100 kbit/s
- Specify the TDM frame structure, assuming that the slot contains *at minimum* $n_i = 8$ bits

Homework

CDM (Code Division Multiplexing)

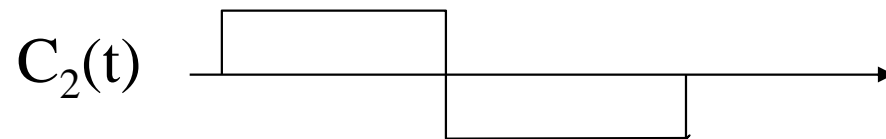
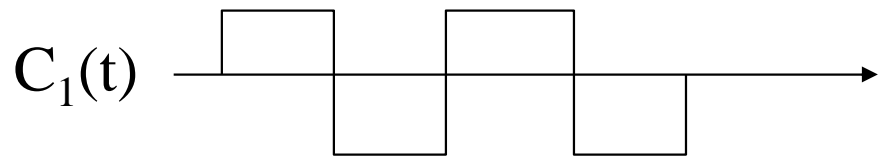
- The CDM technique consists in mixing (i.e., adding) N bit flows (N transmissions), just after having multiplied each one of them with a codeword C_i chosen among the N codewords of an orthogonal code
- Codewords are constituted by N binary symbols, called *chips* in order to distinguish them from *bits*, whose duration is N times shorter than a bit



Orthogonal Codes

- **Orthogonal Signals:** $\int s_1(t) \cdot s_2(t) = 0$

- **Orthogonal Sequencies:**



$$\int_0^T C_1(t) \cdot C_2(t) = 0$$

$$\sum_{i=1}^N c_{1i} \cdot c_{2i} = 0$$

Orthogonal Codes

Hadamard Matrix:

$$H_2 = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$H_{2n} = \begin{bmatrix} H_n & H_n \\ H_n & -H_n \end{bmatrix}$$

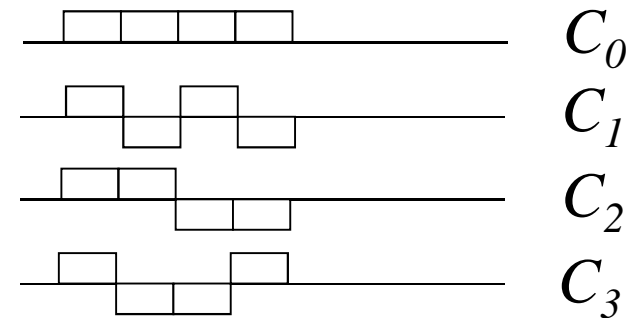
Example for N=4

$$C_0 = \{1, 1, 1, 1\}$$

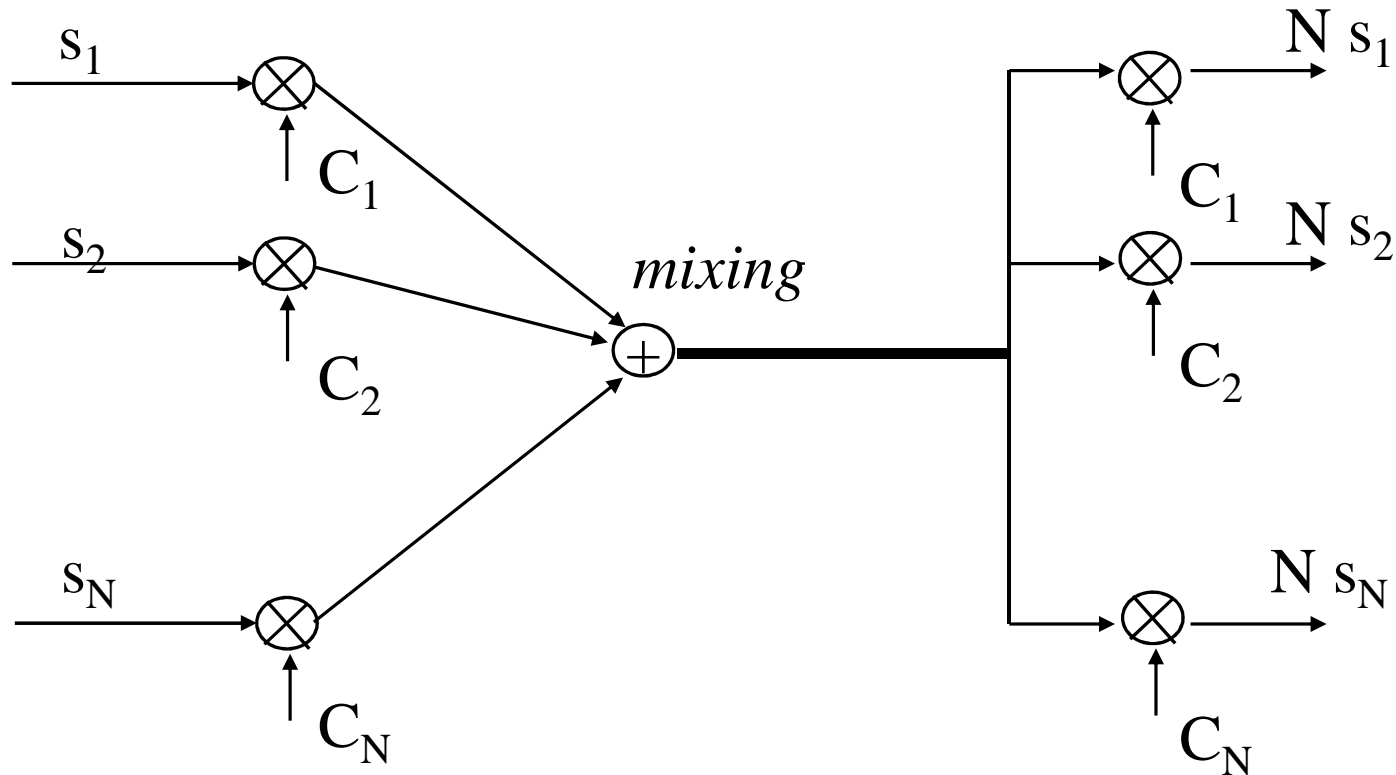
$$C_1 = \{1, -1, 1, -1\}$$

$$C_2 = \{1, 1, -1, -1\}$$

$$C_3 = \{1, -1, -1, 1\}$$



CDM (Code Division Multiplexing)



At the receiver: I can extract the k-th signal by simply multiplying by C_k

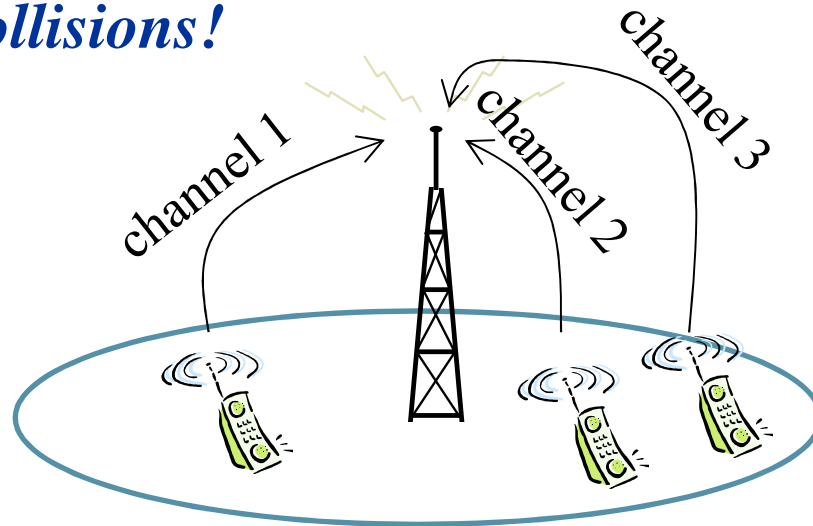
$$\left(\sum_{i=1}^N s_i C_i \right) \cdot C_k = N \cdot s_k$$

WDM (Wavelength Division Multiplexing)

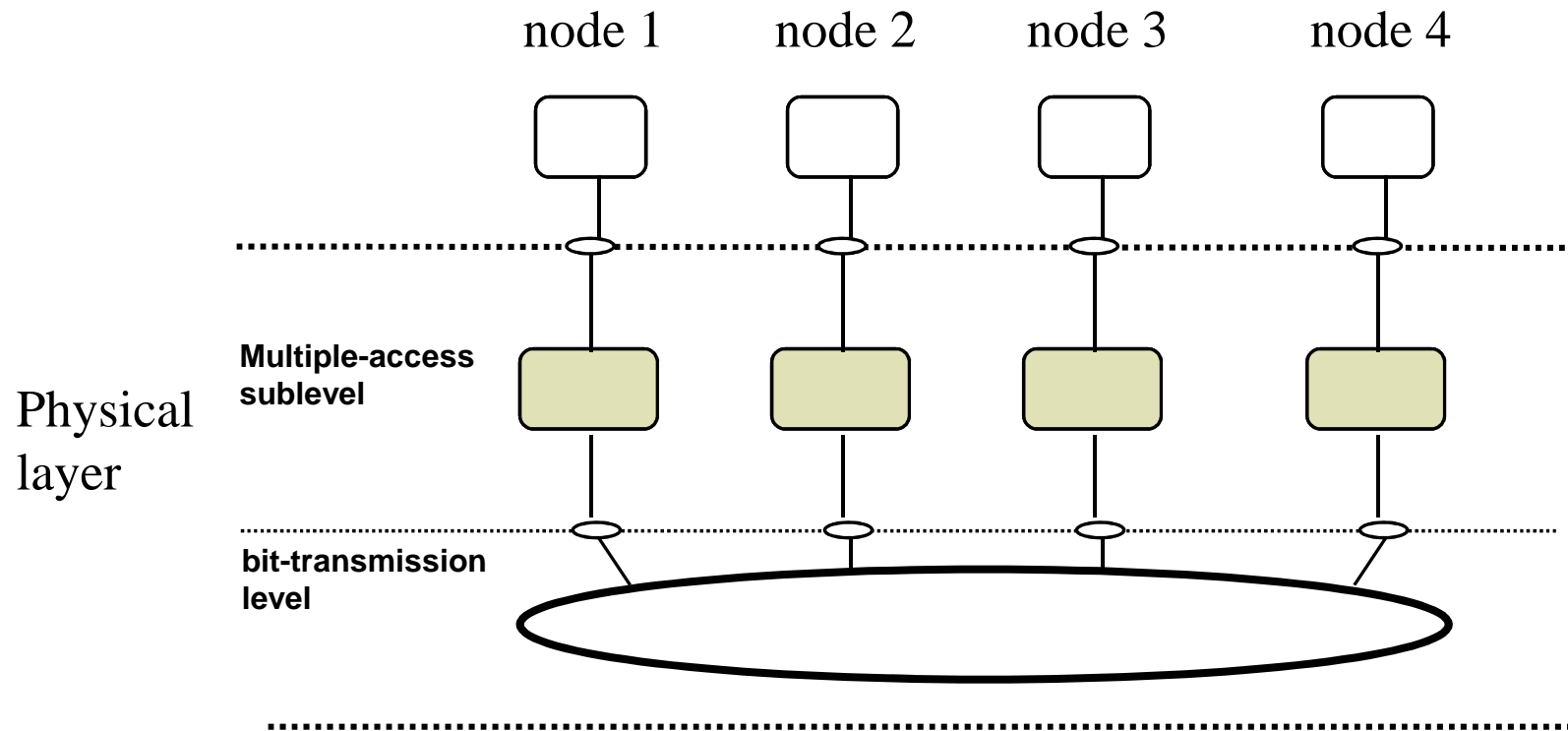
- It's the same as FDM; it is called WDM for historical reasons, related to the development of optical fibers
- Different signals are modulated using different wavelengths on optical fibers
- Each wavelength can carry huge amount of information (5-10 Gbit/s)
- Technological limit: related to the stability of LEDs/Lasers used to modulate signals, as well as by the precision of optical filters
- We have currently commercial devices with 16 – 128 wavelengths (Dense WDM, DWDM)

Multiple Access

- It is *similar* to multiplexing, but conceptually it is *very different*.
- In fact, multiple access is related to *broadcast channels*.
- Hence, the stations/nodes which access the broadcast channel are *distant*, hence they are physically in different places, possibly very far from each other, and so they need to *coordinate among themselves* to access the channel without *collisions!*



Multiple Access: logical scheme



FDMA

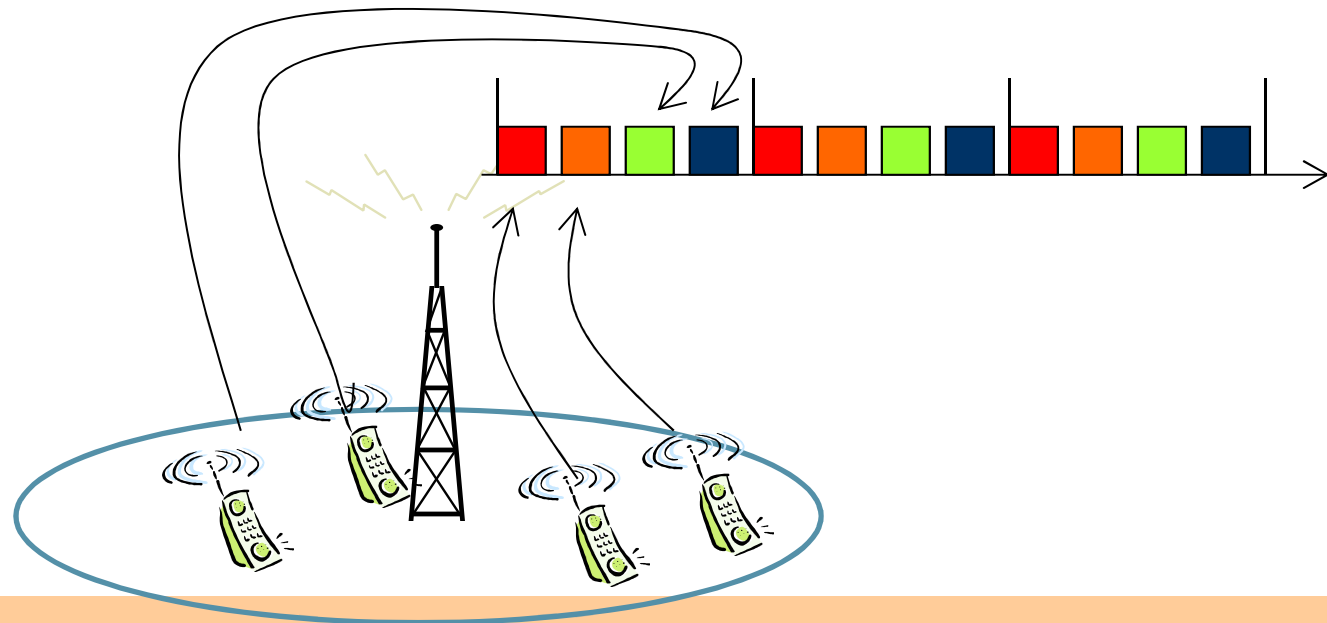
Frequency Division Multiple Access

- **It's analogous to FDM**
- **Different nodes/stations need to coordinate to access the channel, but this is not a problem with FDMA**
- **Examples:**
 - **TV or Radio station broadcast**
 - **Cellular system TACS (Total Access Cellular System) which used 25 kHz subchannels for phone calls**

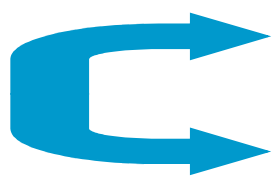
TDMA

Time Division Multiple Access

- It's similar to TDM ...
- ... but here it is necessary for stations to coordinate among themselves to find a common timing reference (necessary to know when slots/frames start and end)
- Synchronization cannot be perfect: *guard times* are necessary to avoid overlapping

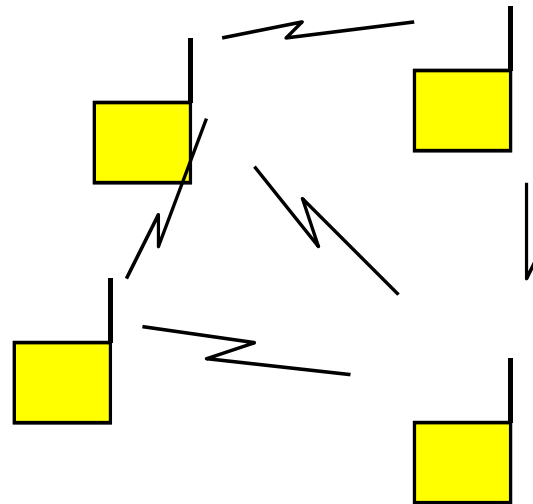
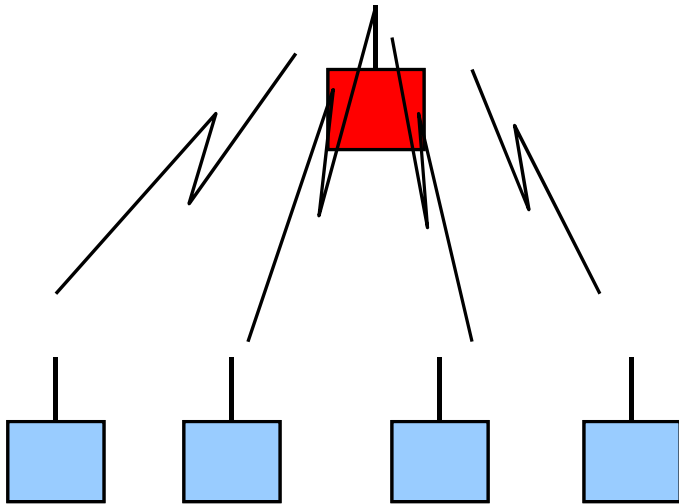


Broadcast channel



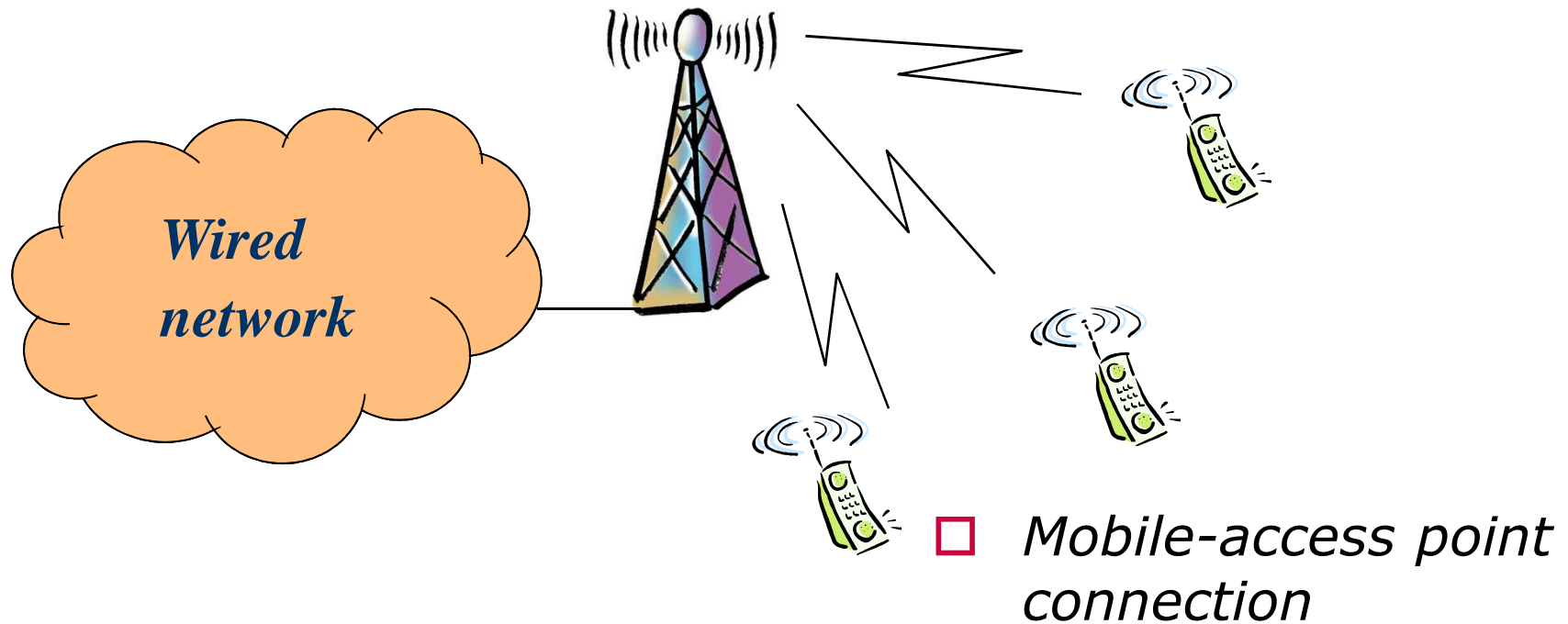
Centralized broadcast channel

Distributed broadcast channel



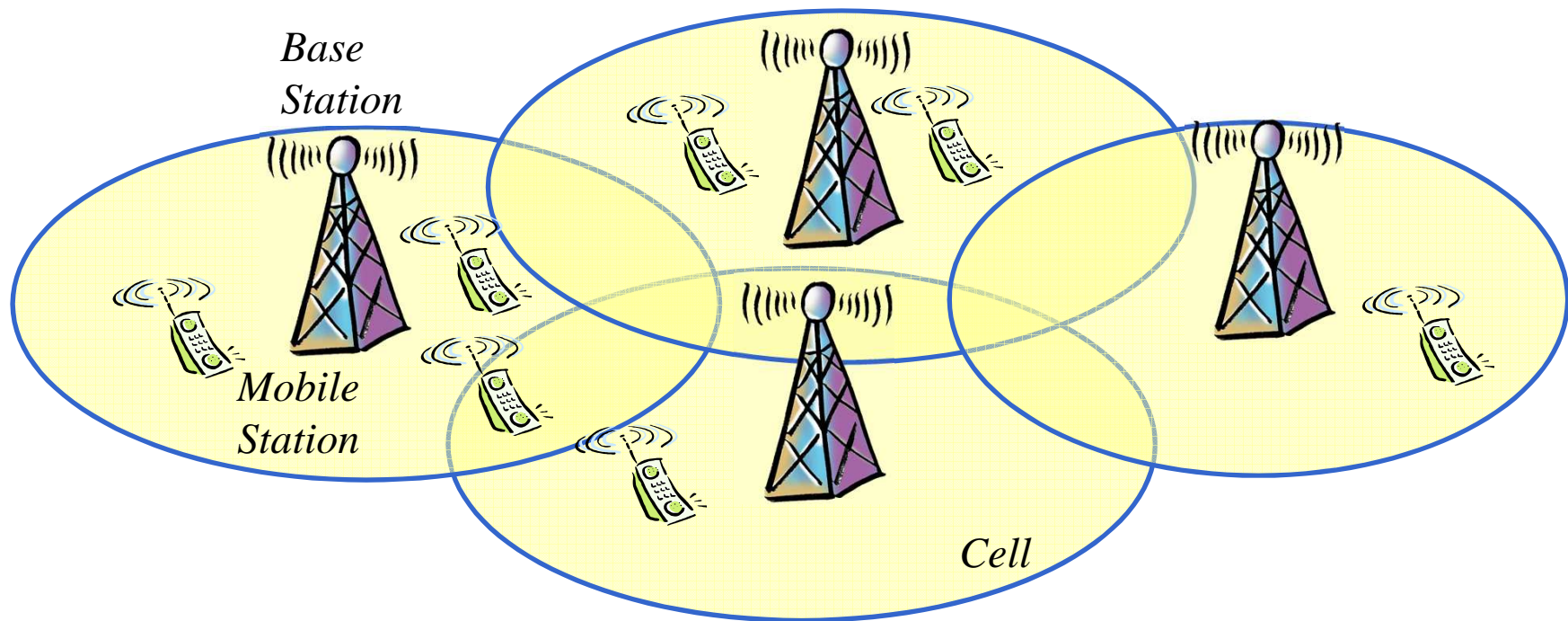
Centralized broadcast channel

- **Fixed access point (cellular systems, WLAN, WMAN)**



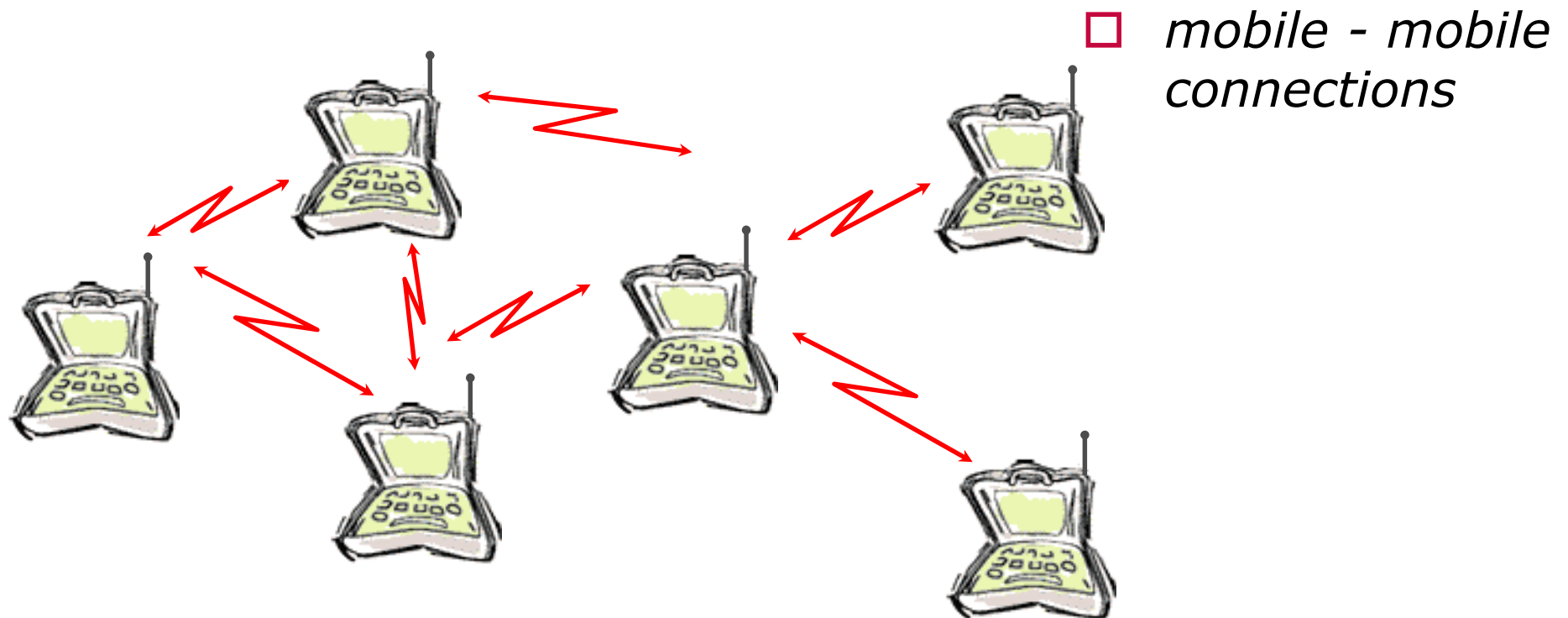
Centralized broadcast channel

- **Cellular coverage:** The territory coverage is obtained by **Base Stations–BS** (or **Access Points**) that provide radio access to **Mobile Stations–MS** within a service area called **CELL**



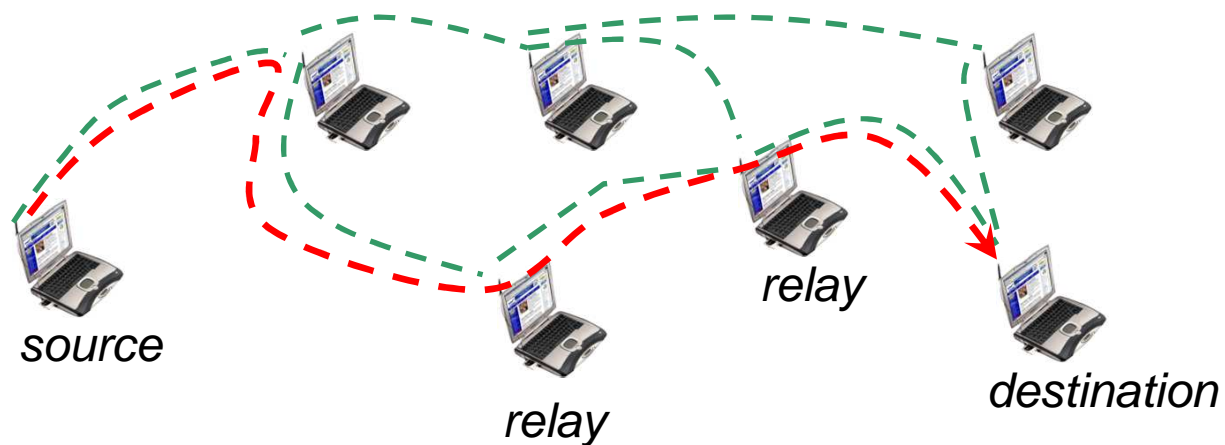
Distributed broadcast channel

- Ad-hoc wireless networks (mesh networks, sensor networks)



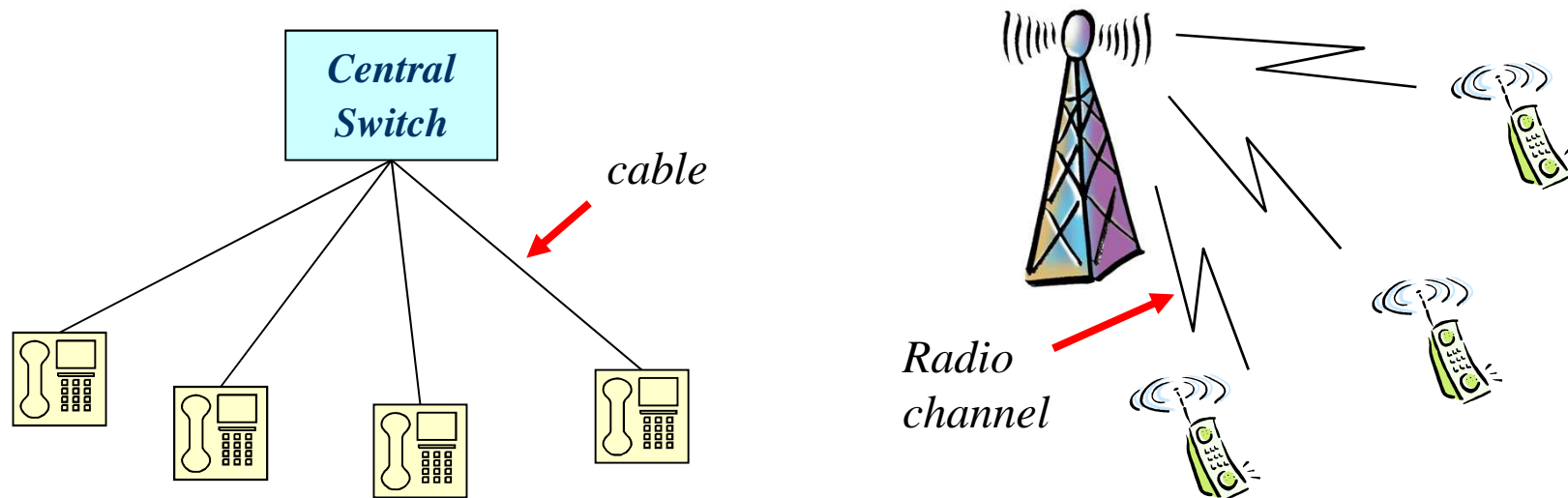
Distributed broadcast channel

- In multi-hop operation mobile stations can forward information



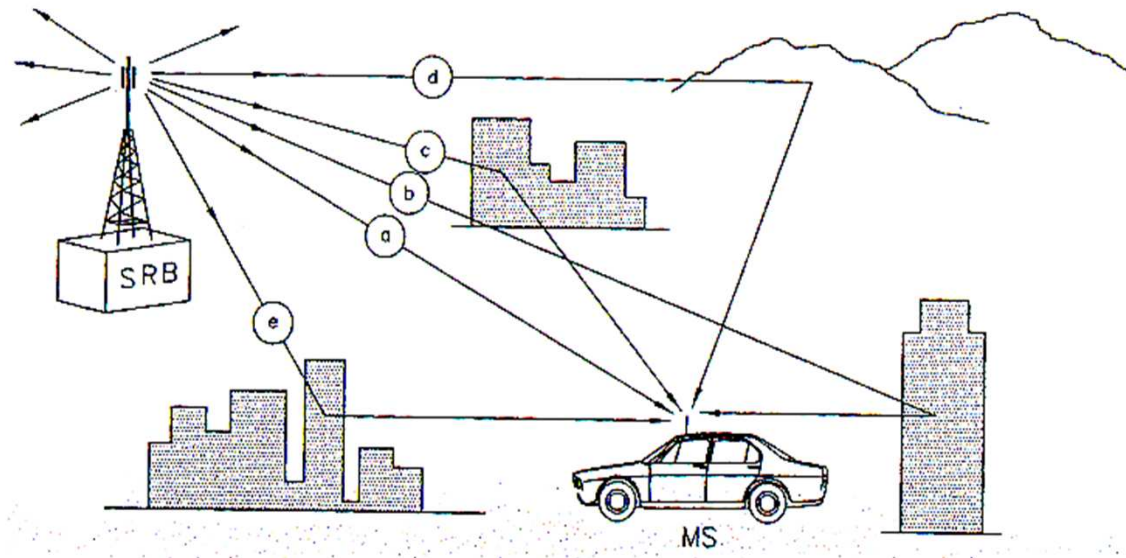
Wired-Wireless networks: Main differences

- **Shared transmission medium**
 - ➔ **Multiple access mechanisms**
 - ➔ **Radio resource reuse**



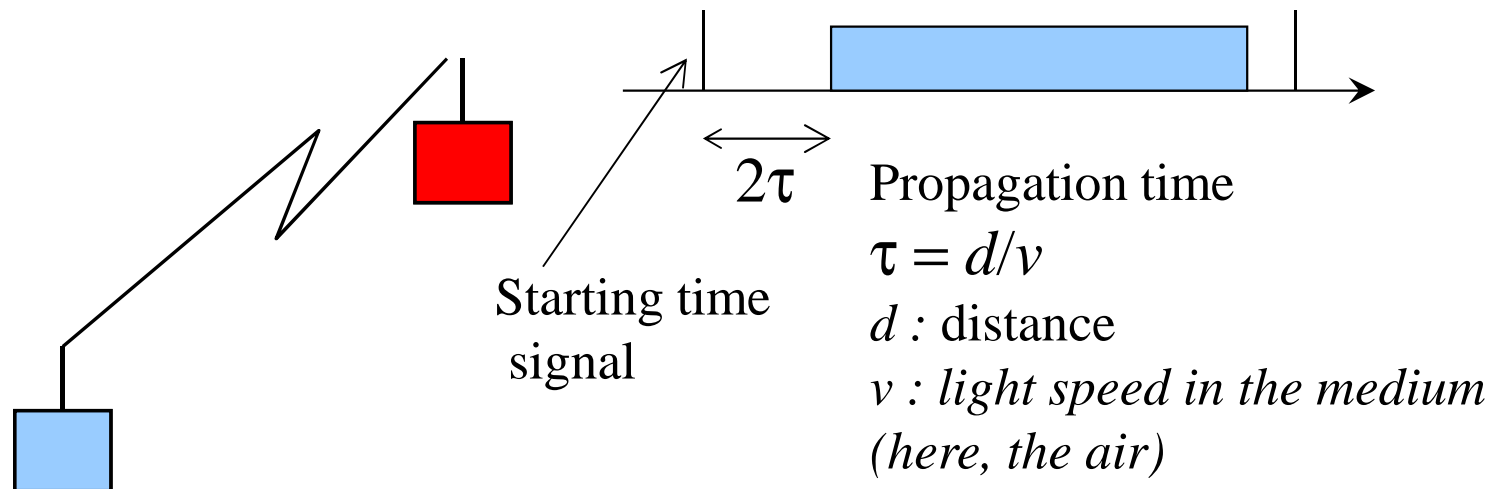
Wired-Wireless networks: Main differences

- **Radio channel**
 - **Variable channel characteristics**
 - **Advanced modulation and coding schemes**



Centralized broadcast channel

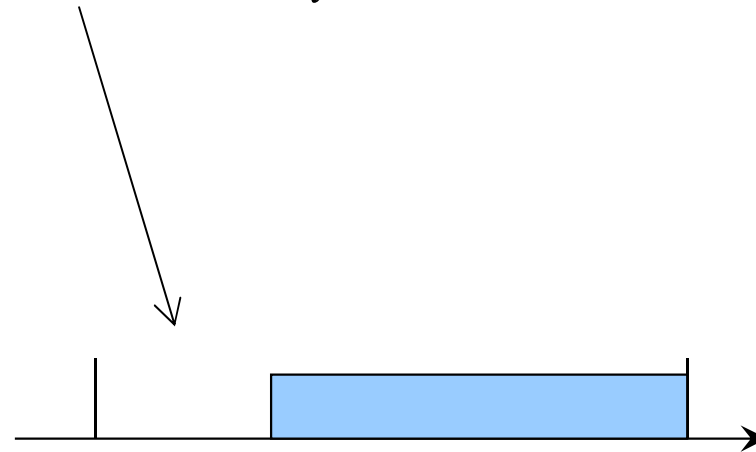
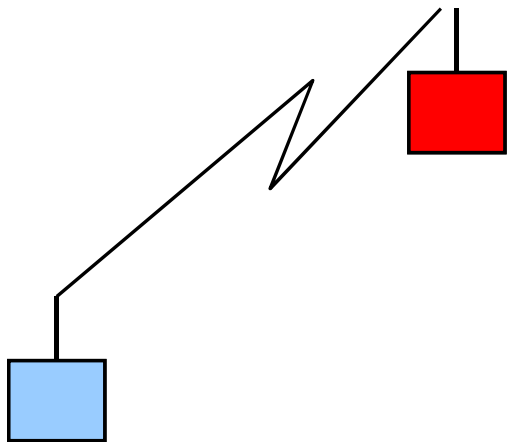
- **The Base Station is vital to enforce synchronization among mobile terminals**
- **Its transmissions are used to synchronize all transmissions (e.g., sending a signal to say when the frame starts)**



Centralized broadcast channel

- Guard time:

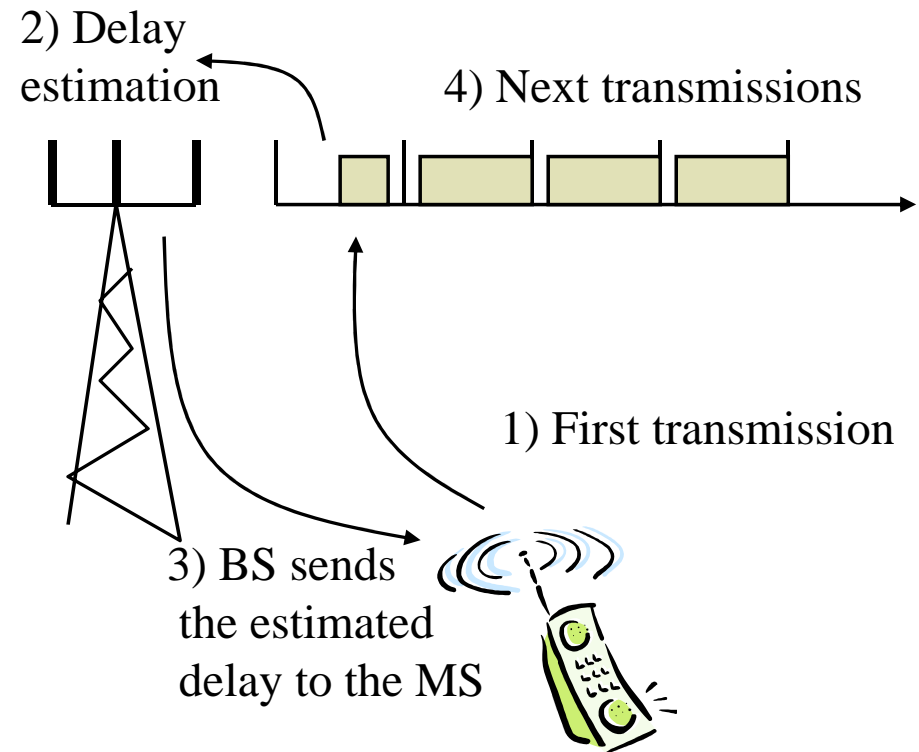
$$T_g = \max_i (2\tau_i)$$



Obviously: the guard time is dominated by the farthest node from the BS

Centralized broadcast channel

- **Timing Advance:**
 - **If each node knows the propagation delay towards the BS, it can anticipate its transmission!**
 - **Propagation delay τ must be estimated (it can be time-varying)**
 - **Estimation error is still possible: time guards are reduced, but they are not null!**
 - **Technique used in GSM**



Efficiency

$$\eta = \frac{T_i}{T_i + T_g} = \frac{1}{1 + \frac{T_g}{T_i}} = \frac{1}{1 + T_g \frac{C}{n_i}}$$

- It depends on the ratio T_g/T_i
- The efficiency decreases:
 - When distances from the BS increase (T_g increases)
 - When the channel speed C increases
 - When the slot duration decreases

CDMA

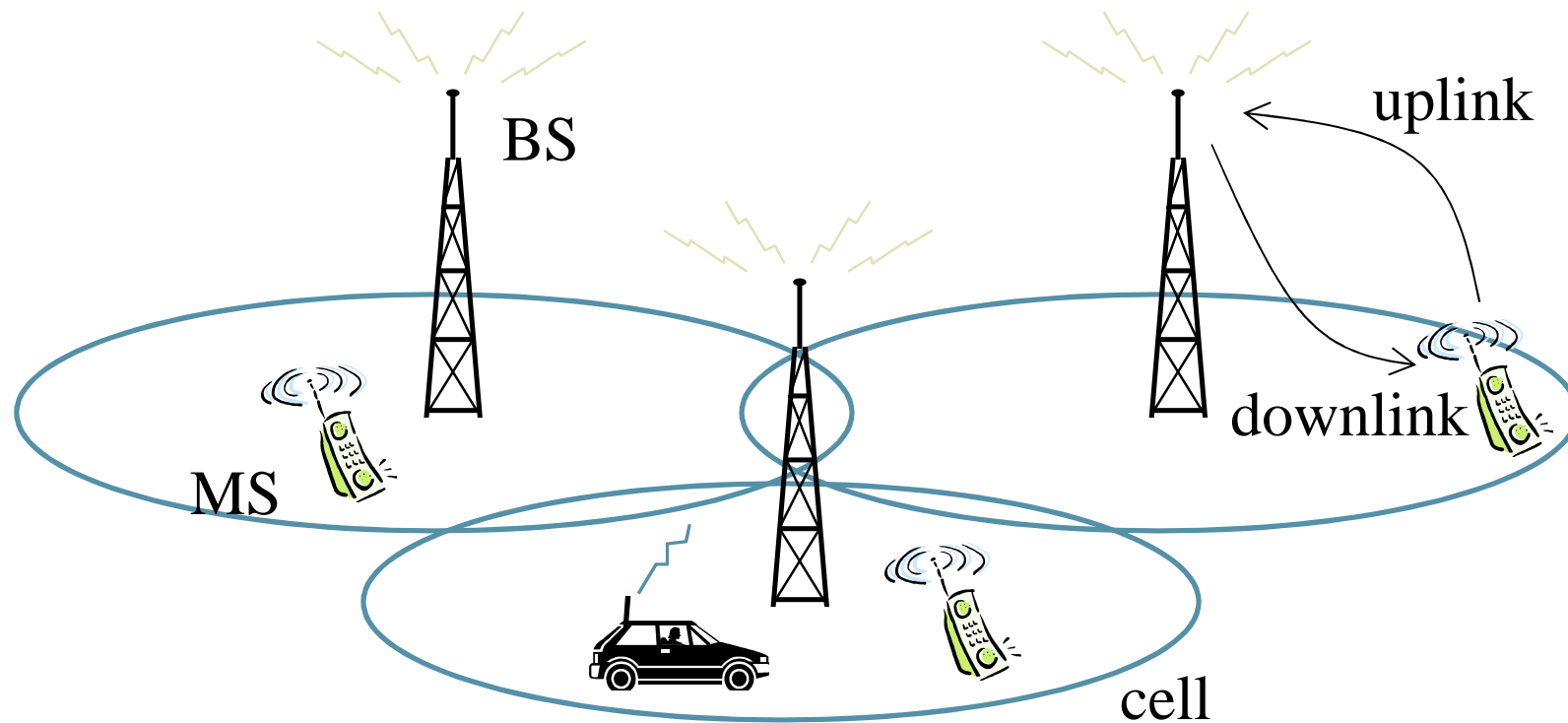
Code Division Multiple Access

- In CDMA it is impossible to have perfect synchronization among different nodes' transmissions
- Hence code orthogonality is *lost*
- We use codes with very low correlation for every possible time shift Δ among themselves
- Used in 3rd generation systems (UMTS)

$$\int_0^T C_1(t) \cdot C_2(t) \neq 0$$

$$\int_0^T C_1(t) \cdot C_2(t+\Delta)$$

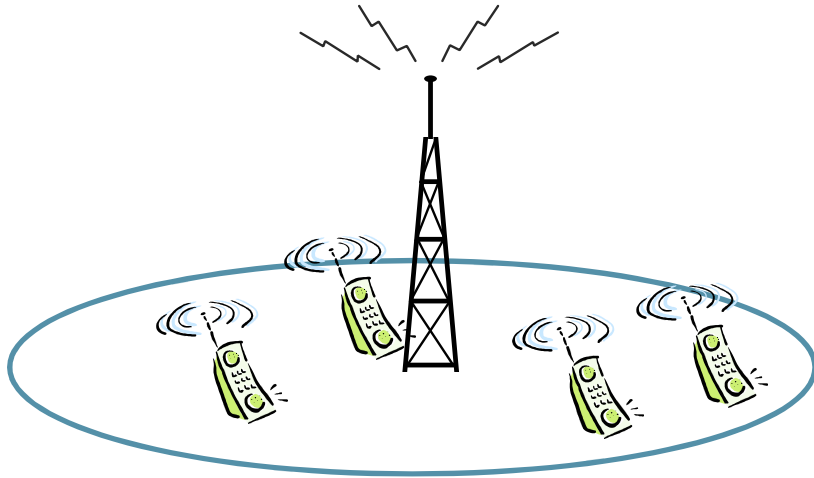
Cellular (Mobile) Systems



MS = Mobile Station
BS = Base Station

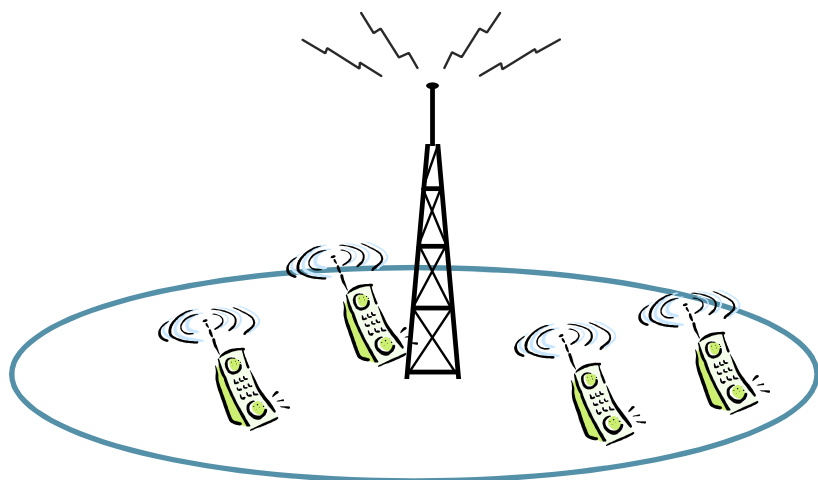
Uplink = from the MS to the BS
Downlink = from the BS to the MS

Radio Access



- The radio access problem is related to the way in which the users in the cell share radio resources
- **downlink:**
 - multiplexing is used
- **uplink:**
 - multiple access is used

Radio Access



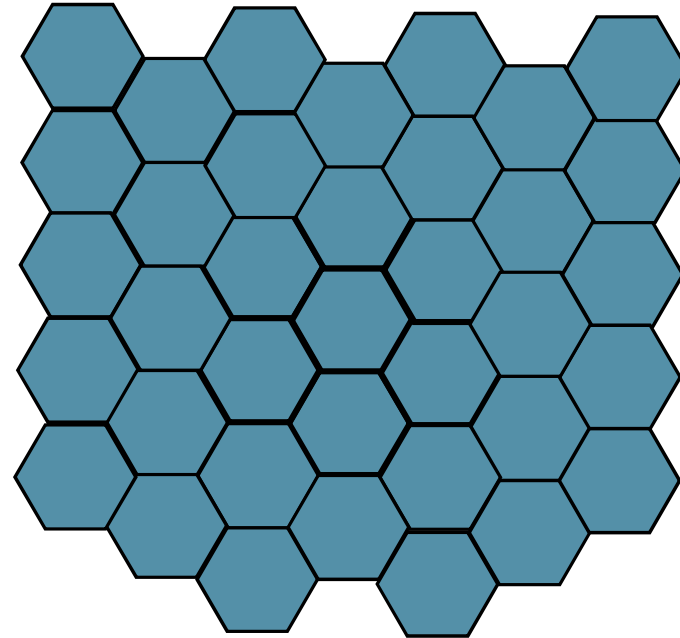
- **First generation systems:**
 - **TACS (Europe)**
AMPS (US)
 - **FDM/FDMA**
(downlink/uplink)
- **Second generation:**
 - GSM (Europe)**
 - D-AMPS (US)**
 - **multi-carrier TDM/TDMA**
- **Third generation:**
 - UMTS**
 - CDM/CDMA**

Frequency reuse

- Available frequencies are not sufficient for all users
- Solution: we reuse the same frequency in different cells (*spatial reuse*)
- Spatial reuse causes *co-channel interference*
- Spatial reuse is made possible if cells are sufficiently far apart so that interference can be small/tolerable (in order to guarantee a good quality of the transmitted signal)

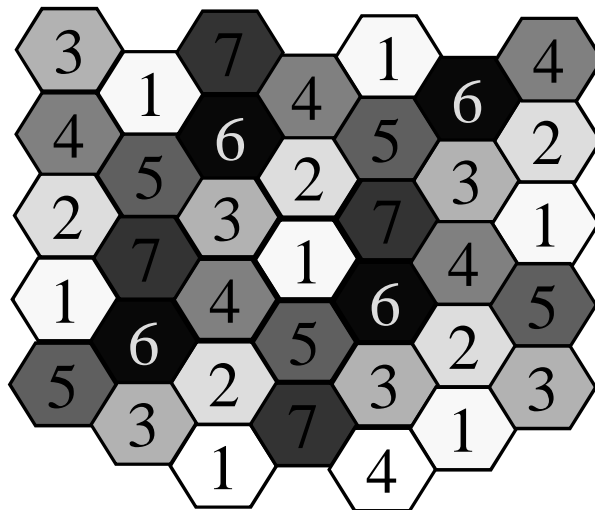
Spatial reuse

- The *interference* is therefore a fundamental, intrinsic feature of cellular systems
- Usually we assume that system quality is good when the ratio between the signal power and the interference power, named *SIR* (Signal-to-Interference Ratio) is higher than a predefined threshold, SIR_{\min}

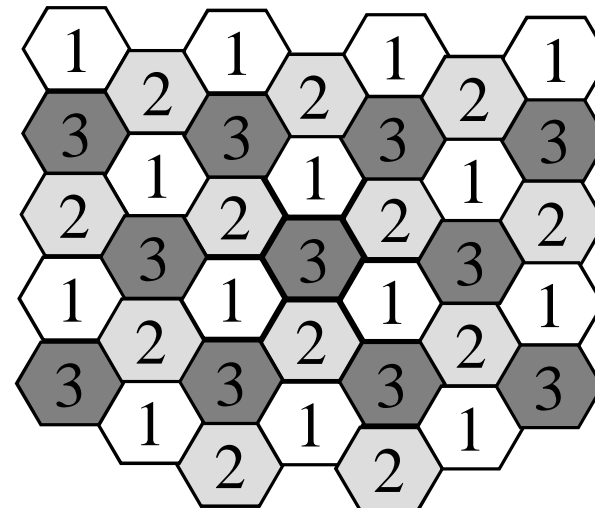


Cluster dimensioning

- All available frequencies are divided into K groups
- We assign a group to each cell in order to maximize the distance between 2 cells that use the same group of frequencies
- Frequency reuse efficiency = $1/K$
- Possible K values: $K=1,3,4,7,9,12,13, \dots$



$K = 7$



$K = 3$

Cluster dimensioning

- If we know/if we set the SIR_{\min} value tolerated by the system, then we can estimate the maximal efficiency of the system, i.e., the minimum K value that can be used
- Received power:

$$P_r = P_t \cdot G \cdot d^{-\eta}$$

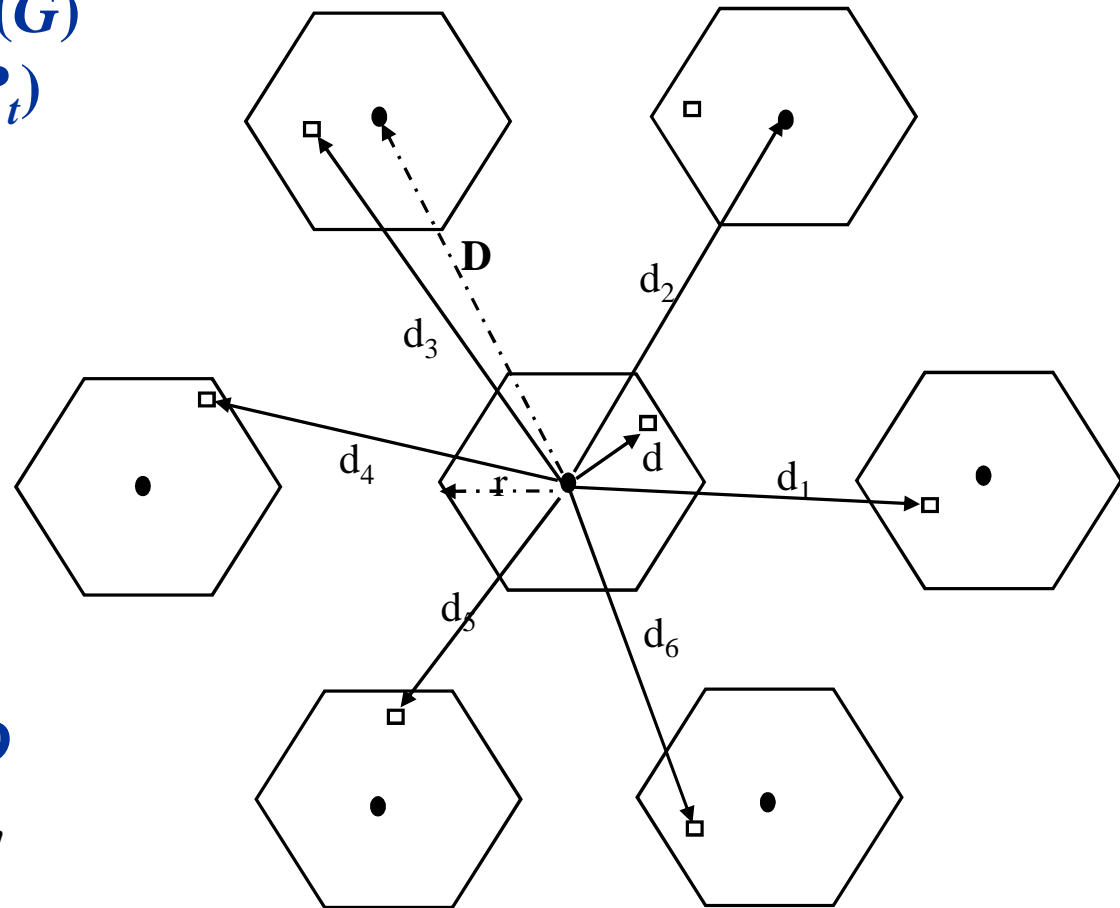
Cluster dimensioning

- Hip.: same antennas (G) and same tx power (P_t)

$$SIR = \frac{P_t \cdot G \cdot d^{-\eta}}{\sum_{i=1}^6 P_t \cdot G \cdot d_i^{-\eta}} = \frac{d^{-\eta}}{\sum_{i=1}^6 d_i^{-\eta}}$$

- Worst case: $d = r$
- Approximation: $d_i = D$

$$SIR \cong \frac{r^{-\eta}}{6D^{-\eta}} = \frac{1}{6} \left(\frac{1}{R} \right)^{-\eta}$$



Cluster dimensioning

- The SIR depends exclusively on the reuse ratio $R=D/r$ (and on η) *but not* on the absolute transmission power or on the cell dimension
- If we fix SIR_{\min} we can compute R_{\min}
- Then, if R_{\min} is known, we can obtain K since we can observe that:

$$K = \frac{R^2}{3}$$

- and therefore:

$$K_{\min} = \frac{(6 SIR)^{2/\eta}}{3}$$

Exercise

- Let us dimension a cluster for a cellular system that tolerates $SIR_{\min} = 18$ dB, considering the case where the path-loss exponent η is equal to 3.9

$$K_{\min} = \frac{(6 SIR)^{2/\eta}}{3} = \frac{(6 \cdot 63.1)^{2/3.9}}{3} = 6.99$$

Summary

- **dB**
 - **Logarithmic scale**
 - **If we use absolute powers**

$$P_{dB} = 10 \log_{10} P$$

$$P = 10^{P_{dB}/10}$$

Summary

- The *product* in linear scale corresponds to a *sum* using dB
- The *ratio* corresponds to a *difference* in dB

$$G \cdot P \rightarrow G_{dB} + P_{dB}$$

$$P / A \rightarrow P_{dB} - A_{dB}$$

Summary

- **Notable values**

$$2 \rightarrow 3dB$$

$$3 \rightarrow 4.77dB$$

$$4 = 2 \cdot 2 \rightarrow 3 + 3 = dB$$

$$5 \rightarrow 7dB$$

$$6 \rightarrow 7.77dB$$

$$8 \rightarrow 9dB$$

$$9 \rightarrow 9.54dB$$

$$10 \rightarrow 10dB$$

$$100 \rightarrow 20dB$$

$$1000 \rightarrow 30dB$$