Année 2022-23

Prof. Burkhart Wolff wolff@lri.fr
www.lri.fr/~wolff/teach-material/2022-2023/M2-CSMR/

## **TP 5 - Modeling Operational Semantics**

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Objective : Defining an operational semantics "Plotking-style" big-step semantic for as more-or-less standard regular expression language. This type of semantics represents "states" 's in the original language and inductively models the transition relation between states via a transition predicate  $\langle \_, \_ \rangle \longrightarrow_c \_$  of type 's ×' output option ×' s  $\Rightarrow$  bool where option is the usual option type constructor from Main. Following common terminology in automata theory, we will call a list of atoms a word, and a set of words a language. The denotational semantics we are referring here is a function that maps an (abstract) syntax to a set of deonotations, i.e. regular expressions to the language they denote. We will introduce the term epsilon as abbreviation for Star Empty, an  $\alpha$  rexp term for the empty language.

We reuse the abstract syntax  $\alpha$  rexp of the regular expression language from TP4.

## Exercice 1 (Inductive Sets - Inductive Proofs)

Define a Plotkin-style semantics for regular expressions, where 'output is set to ' $\alpha$ . Complete the list of inductive rules starting with :

- $\langle \lfloor a \rfloor, Some \ a \rangle \longrightarrow_c epsilon$
- $\langle Empty : R, None \rangle \longrightarrow_c Empty$
- $-\langle |a|: R, Some |a\rangle \longrightarrow_c R$
- $\langle Star \ r, None \rangle \longrightarrow_c (r : (Star \ r))$
- $-\langle Star \ r, None \rangle \longrightarrow_{c} epsilon$

— ... Tasks :

- 1. Prove  $\langle \lfloor a \rfloor$ , Some  $a' \rangle \longrightarrow_c epsilon = (a' = a)$  and  $\langle \lfloor a \rfloor, a' \rangle \longrightarrow_c R' = (a' = Some \ a \land R' = epsilon)$  (This should hold for your completion of the above inductive rule-set).
- 2. Derive all similar lemmas resulting from your definitions (should be approx 8). Hint : for the latter rule, there is actually a specific command that derives this type of simplification lemmas automatically. For example, the last mentioned lemma could be derived automatically via :

inductive simps atom 1S : "
$$\langle |a|, a' \rangle \longrightarrow_c R'$$
"

3. Prove the eliminiation rule :

$$\langle \lfloor a \rfloor, Some \ a' \rangle \longrightarrow_c epsilon \Longrightarrow (a' = a \Longrightarrow P) \Longrightarrow P$$

4. Prove all other elimination rules and configure them into the global context as such. Hint : for the latter rule, there is actually a specific command that derives this type of simplification lemmas automatically. For example, the last mentioned lemma could be derived automatically via :

inductive simps atom 1S : " $\langle |a|, a' \rangle \longrightarrow_c R'$ "

5. Now define the mu;tiple step semantics Plotkin style. This hould lopok like this : inductive

| evalstar ::               | "['a rexp,'a list,'a rexp] $\Rightarrow$ bool" (" $\langle , \rangle / \longrightarrow_c^* \_$ " [0,0,60] 60)   |
|---------------------------|---|
| where                     |   |
| idle:                     | "(epsilon,[]) $\longrightarrow_{c}^{*}$ epsilon"  |
| step1:                    | $\langle r,Some  a \rangle \longrightarrow_{c} r' \Longrightarrow \langle r,[a] \rangle \longrightarrow_{c}^{*} r''$  |
| <pre>continuation1:</pre> | $\langle r,None\rangle \longrightarrow_{c} r' \implies \langle r',S\rangle \longrightarrow_{c}^{*} r'' \implies \langle r,S\rangle \longrightarrow_{c}^{*} r'''$                    |
| continuation2:            | $\langle r,Some a\rangle \longrightarrow_{c} r' \Longrightarrow \langle r',S\rangle \longrightarrow_{c}^{*} r'' \Longrightarrow \langle r,a\#S\rangle \longrightarrow_{c}^{*} r'''$ |

6. Prove :

$$\langle Star((\lfloor CHR''a''\rfloor | \lfloor CHR''b''\rfloor) : \lfloor CHR''c''\rfloor), ''bc''\rangle \longrightarrow_c^* epsilon$$

7. Prove :

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theorem operational_implies_denotational_generalized':
assumes nat_steps: "(⟨r,s⟩ → c* r')"
and den_cont: "∃as. as ∈ L r'"
shows "∃ as . s@as ∈ L r ∧ as ∈ L r'"
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8. Prove the main theorem "operational semantics implies denotational semantics" :

$$(\langle r, s \rangle \longrightarrow_c^* epsilon) \longrightarrow s \in L(r)$$

Note : Main provides the notation CHR ''a'' for "the character a". Strings are defined as lists of characters.

## Exercice 2 (OPTIONAL : Report )

(IN CASE THAT YOU WANT TO HAVE IT GRADED. RECALL THAT 2 OUT OF 6 TP's SHOULD BE SUBMITTED.)

1. Write a little report answering all questions above, note the difficulties you met, add some screenshots if appropriate. 3 pages max (except screenshots and other figures).