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https://www.lri.fr/~wolff/teach-material/2023-2024/M2-CSMR/index.html

TP 5 - Modeling Operational Semantics

Semaine du 6 fevrier 2024

Objective: Defining an operational semantics "Plotking-style" big-step semantic for as more-or-less standard regular expression language. This type of semantics represents "states" 's in the original language and inductively models the transition relation between states via a transition predicate $\langle _, _ \rangle \longrightarrow_c _$ of type 's ×' output option ×' s \Rightarrow bool where option is the usual option type constructor from Main. Following common terminology in automata theory, we will call a list of atoms a word, and a set of words a language. The denotational semantics we are referring here is a function that maps an (abstract) syntax to a set of deonotations, i.e. regular expressions to the language they denote. We will introduce the term epsilon as abbreviation for Star Empty, an α rexp term for the empty language.

We reuse the abstract syntax α rexp of the regular expression language from TP4.

Exercice 1 (Inductive Sets - Inductive Proofs)

Define a Plotkin-style semantics for regular expressions, where 'output is set to ' α . Complete the list of inductive rules starting with :

- $-\langle |a|, Some \ a \rangle \longrightarrow_c epsilon$
- $--\langle Empty: R, None \rangle \longrightarrow_c Empty$
- $-\langle |a|: R, Some\ a\rangle \longrightarrow_c R$
- $\langle Star \ r, None \rangle \longrightarrow_c (r : (Star \ r))$
- $\langle Star \ r, None \rangle \longrightarrow_c epsilon$
- ...

Tasks:

- 1. Prove $\langle \lfloor a \rfloor, Some \ a' \rangle \longrightarrow_c epsilon = (a' = a)$ and $\langle \lfloor a \rfloor, a' \rangle \longrightarrow_c R' = (a' = Some \ a \land R' = epsilon)$ (This should hold for your completion of the above inductive rule-set).
- 2. Derive all similar lemmas resulting from your definitions (should be approx 8). Hint: for the latter rule, there is actually a specific command that derives this type of simplification lemmas automatically. For example, the last mentioned lemma could be derived automatically via:

inductive simps atom1S: "
$$\langle |a|, a' \rangle \longrightarrow_c R'$$
"

3. Prove the eliminiation rule:

$$\langle |a|, Some \ a' \rangle \longrightarrow_c epsilon \Longrightarrow (a' = a \Longrightarrow P) \Longrightarrow P$$

4. Prove all other elimination rules and configure them into the global context as such. Hint: for the latter rule, there is actually a specific command that derives this type of simplification lemmas automatically. For example, the last mentioned lemma could be derived automatically via:

$$inductive_simps\ atom1S\ :\ "\langle |a|,a'\rangle \longrightarrow_c R'"$$

5. Now define the mu; tiple step semantics Plotkin style. This hould lopok like this:

6. Prove:

$$\langle Star((\lfloor CHR''a''\rfloor | \lfloor CHR''b'' \rfloor) : \lfloor CHR''c'' \rfloor), "bc'' \rangle \longrightarrow_c^* epsilon$$

7. Prove:

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theorem operational_implies_denotational_generalized': assumes nat_steps: "(\langle r,s\rangle \longrightarrow_c^* r')" and den_cont: "\existsas. as \in L r'" shows "\exists as . s@as \in L r \land as \in L r'"
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8. Prove the main theorem "operational semantics implies denotational semantics":

$$(\langle r, s \rangle \longrightarrow_c^* epsilon) \longrightarrow s \in L(r)$$

Note: Main provides the notation CHR ''a' for "the character a". Strings are defined as lists of characters.

Exercice 2 (OPTIONAL : Report)

(IN CASE THAT YOU WANT TO HAVE IT GRADED. RECALL THAT 2 OUT OF 6 TP's SHOULD BE SUBMITTED.)

1. Write a little report answering all questions above, note the difficulties you met, add some screenshots if appropriate. 3 pages max (except screenshots and other figures).