

FROM RESEARCH TO INDUSTRY

cea tech

list

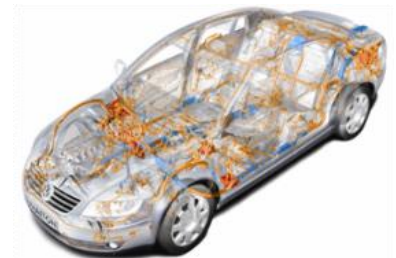
COMBINING FORMAL TOOLS

SOURCE CODE VERIFICATION OF FUNCTIONAL TEMPORAL PROPERTIES

L. CORRENSON, P. BAUDIN

F. KIRCHNER ★

CONTEXT | CRITICAL COMMAND AND CONTROL SOFTWARE



CONTEXT | DEVELOPMENT CYCLE

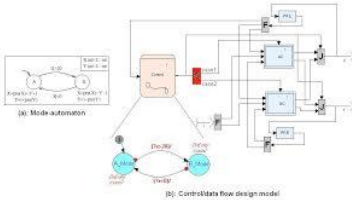
Design



Validation



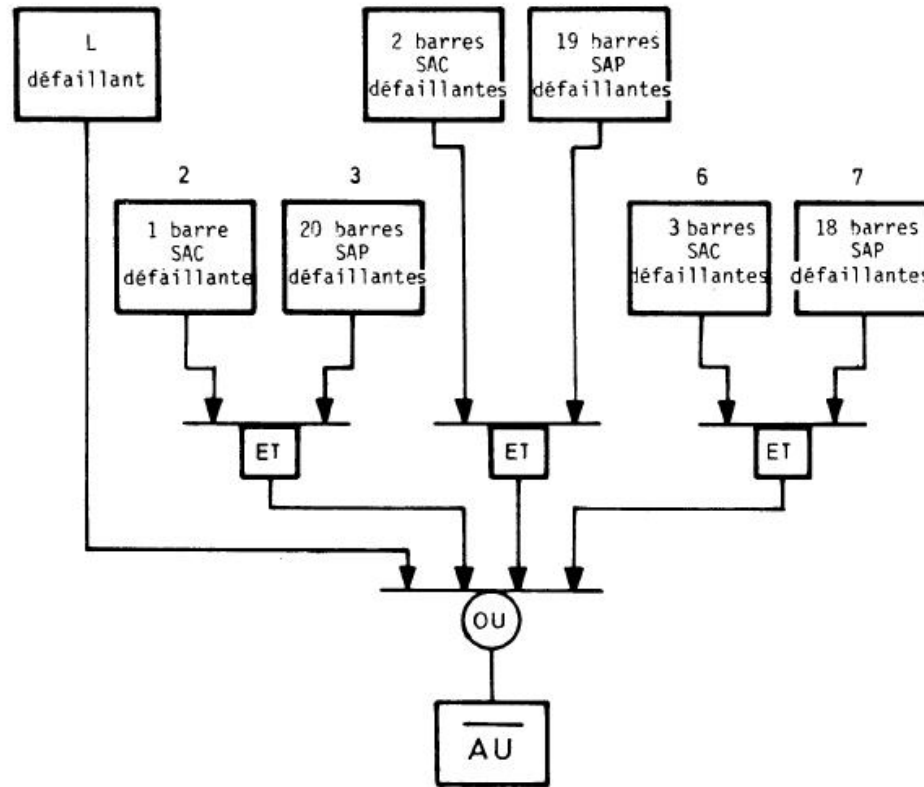
```
int div(int a, int b) {
  int q;
  q=0;
  while (a
    a=a-b;
    q=q+1;
  }
  return q
}
```



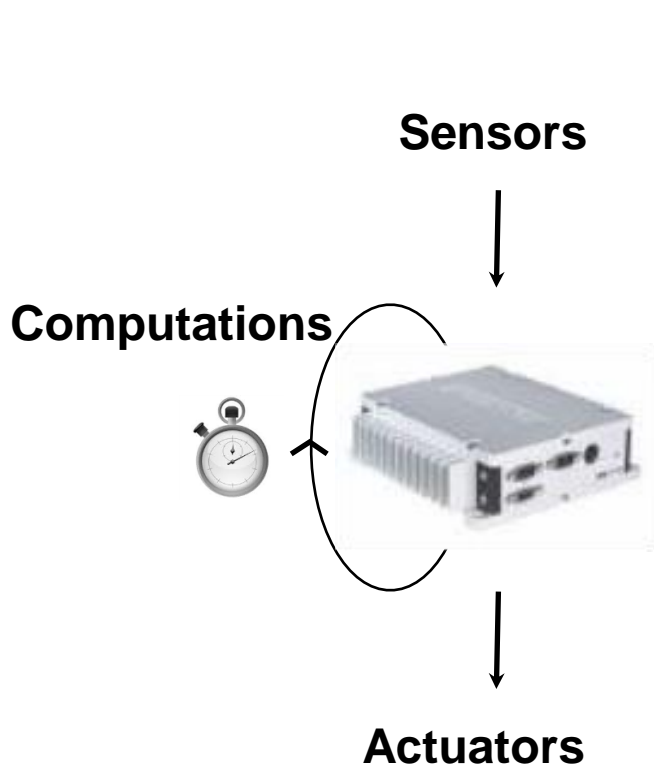
Implementation



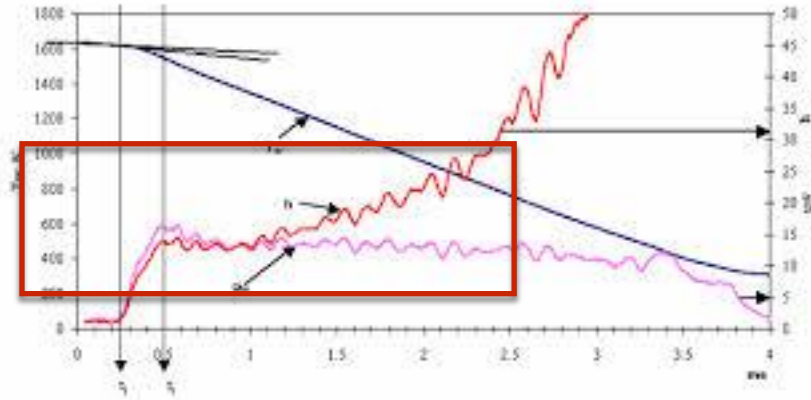
DESIGN | EXPECTED PROPERTIES



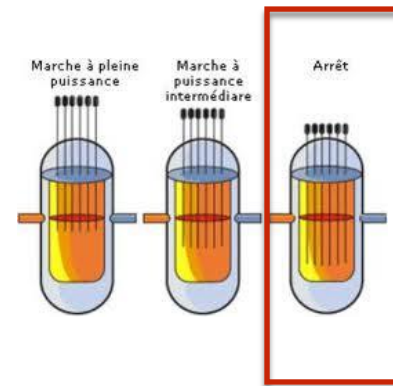
DESIGN | EXPECTED PROPERTIES



ASSUME



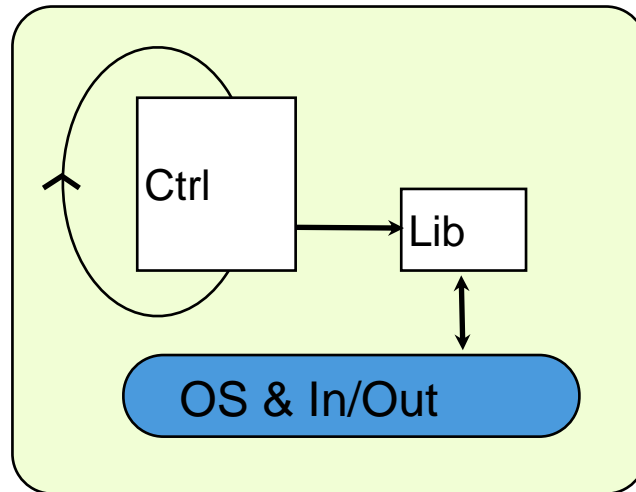
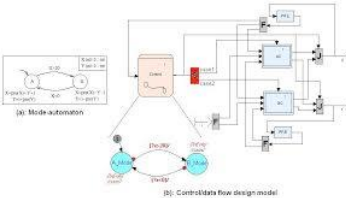
ASSERT



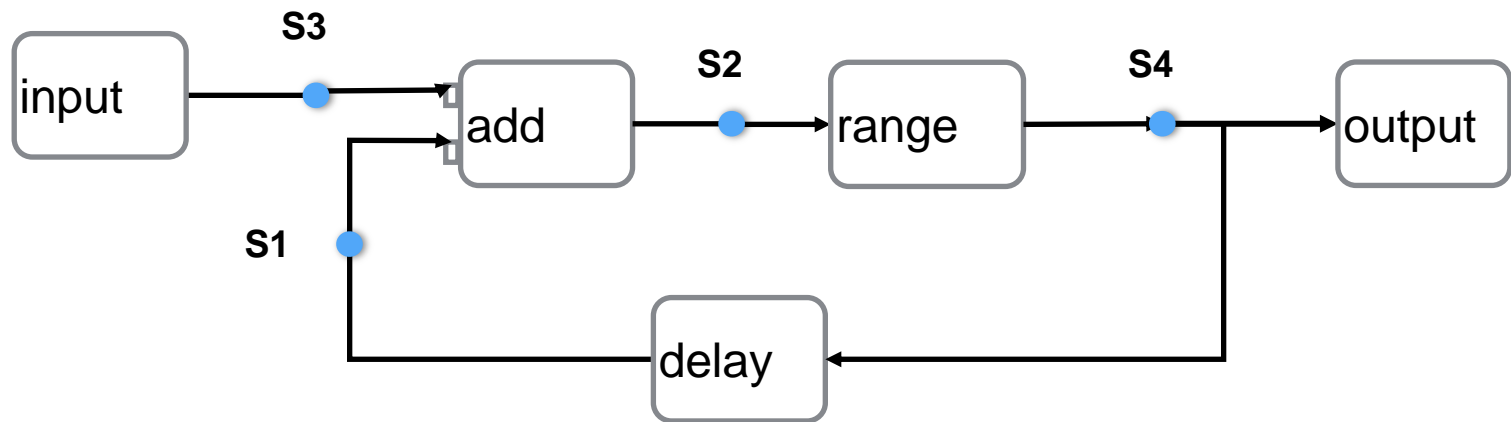
DESIGN | WHAT KIND OF CODE?



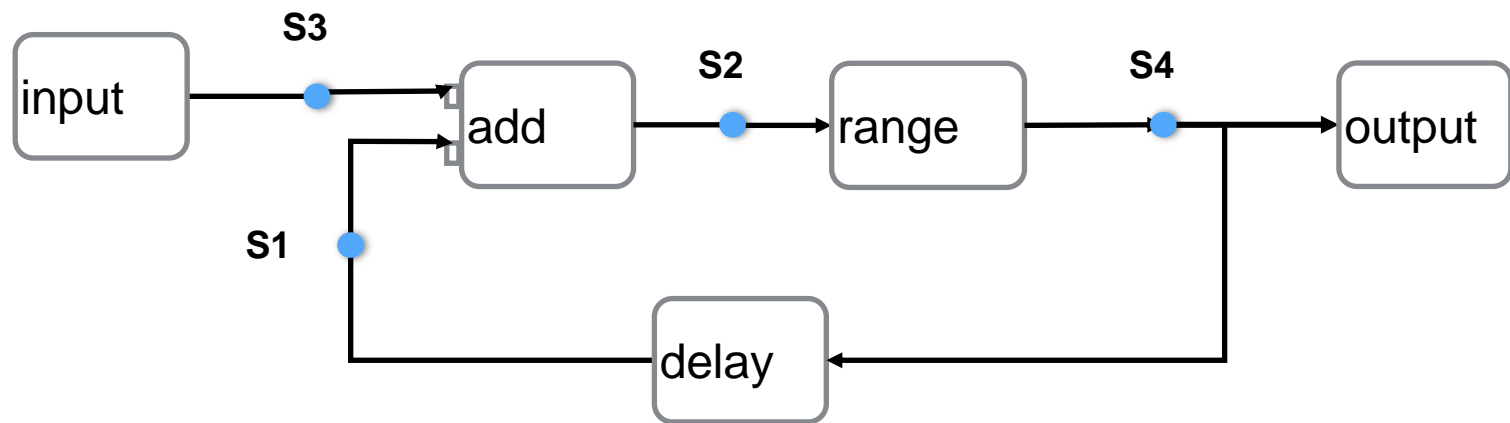
```
int div(int a,
int q;
q=0;
while (a>=b)
a=a-b;
q=q+1;
}
return q;
}
```



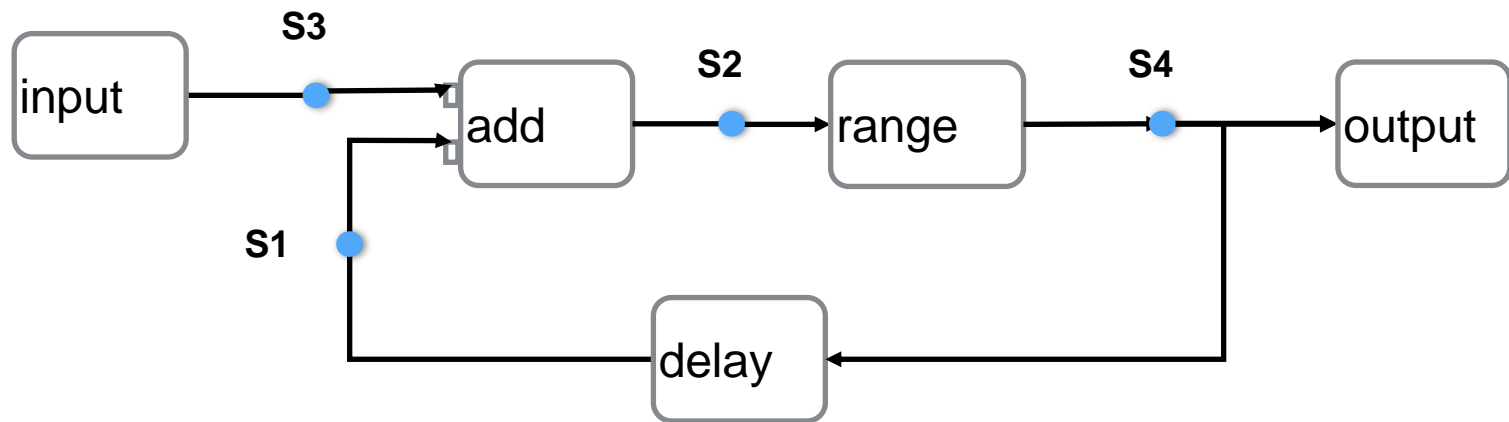
DESIGN | DIAGRAMS



DESIGN | DIAGRAMS: AN ENGINEER'S CODE...



DESIGN | ... IS A COMPUTER SCIENTIST'S MODEL



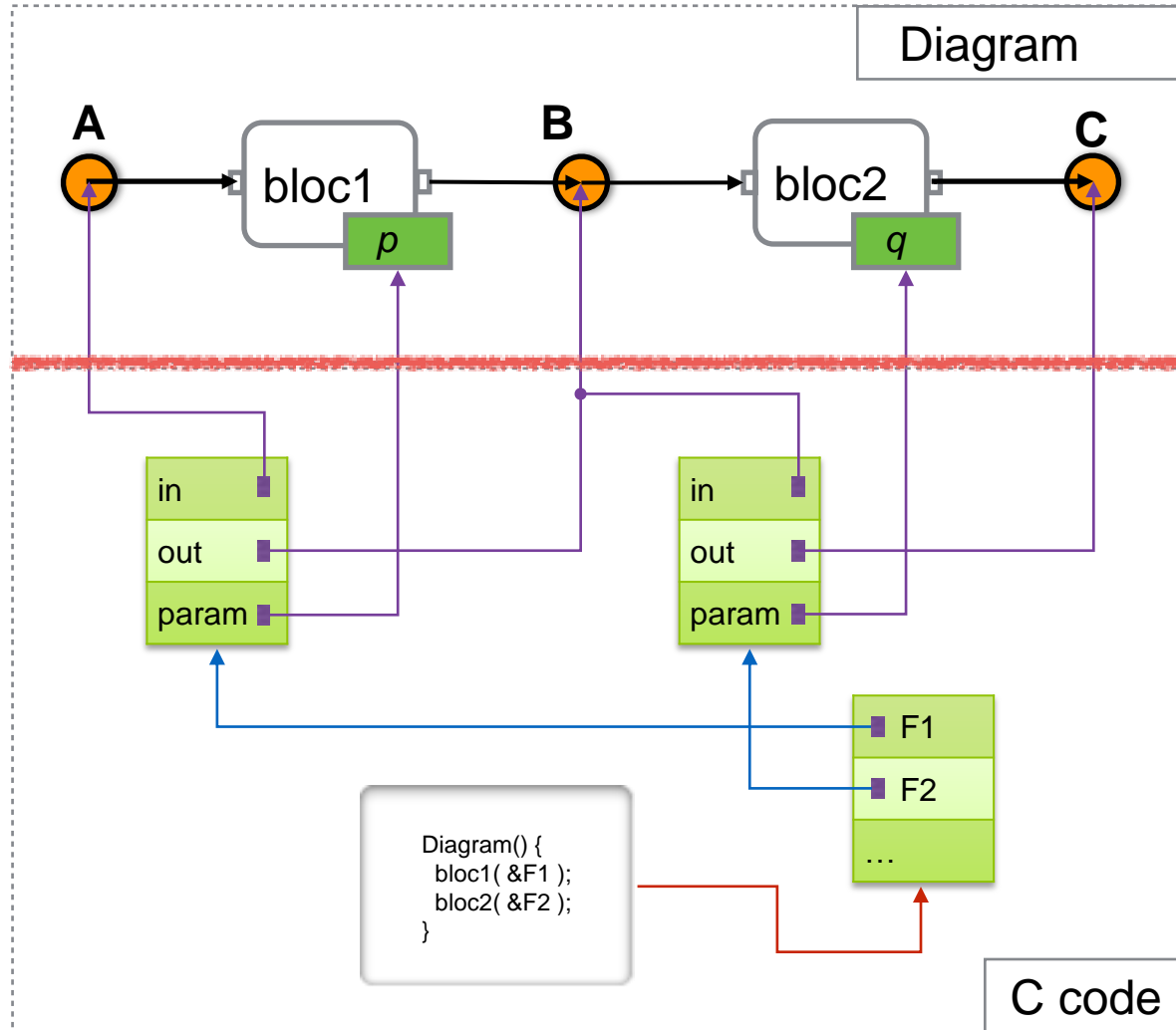
DESIGN | DIAGRAM INVOCATION

```

main()
{
    mode = Init ;
    Diagram(); // Initialization phase
    mode = Param ;
    Diagram(); // Parametrization phase ———  $\varphi$ 
    mode = Run ;
    loop {
        input(); // Sensor reads
        Diagram(); // Computational cycle ———  $M$ 
        output(); // Actuator writes
    }
}

```

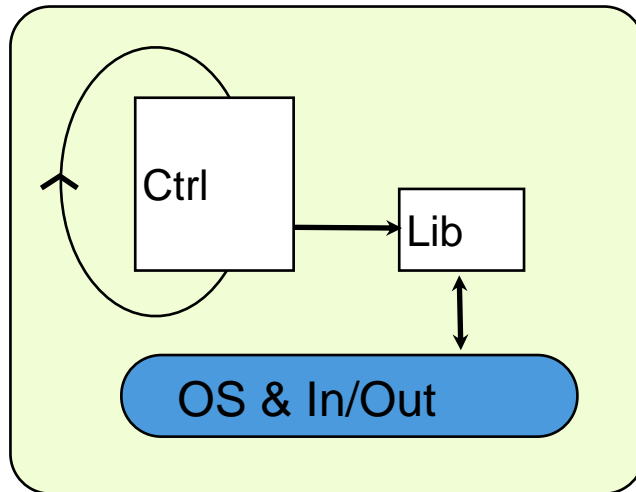
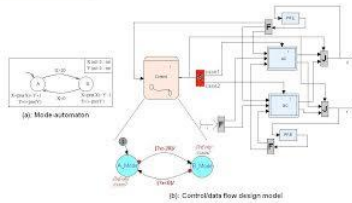
DESIGN | DIAGRAM DECOMPOSITION



WHAT KIND OF VERIFICATION?



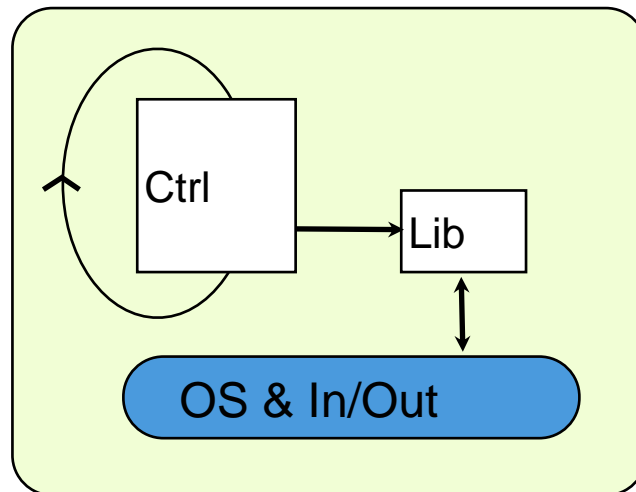
```
int div(int a, int b) {
  int q;
  q=0;
  while (a>=b)
    a=a-b;
    q=q+1;
  }
  return q;
}
```



CAN WE ATTEMPT MODEL-LESS VERIFICATION?

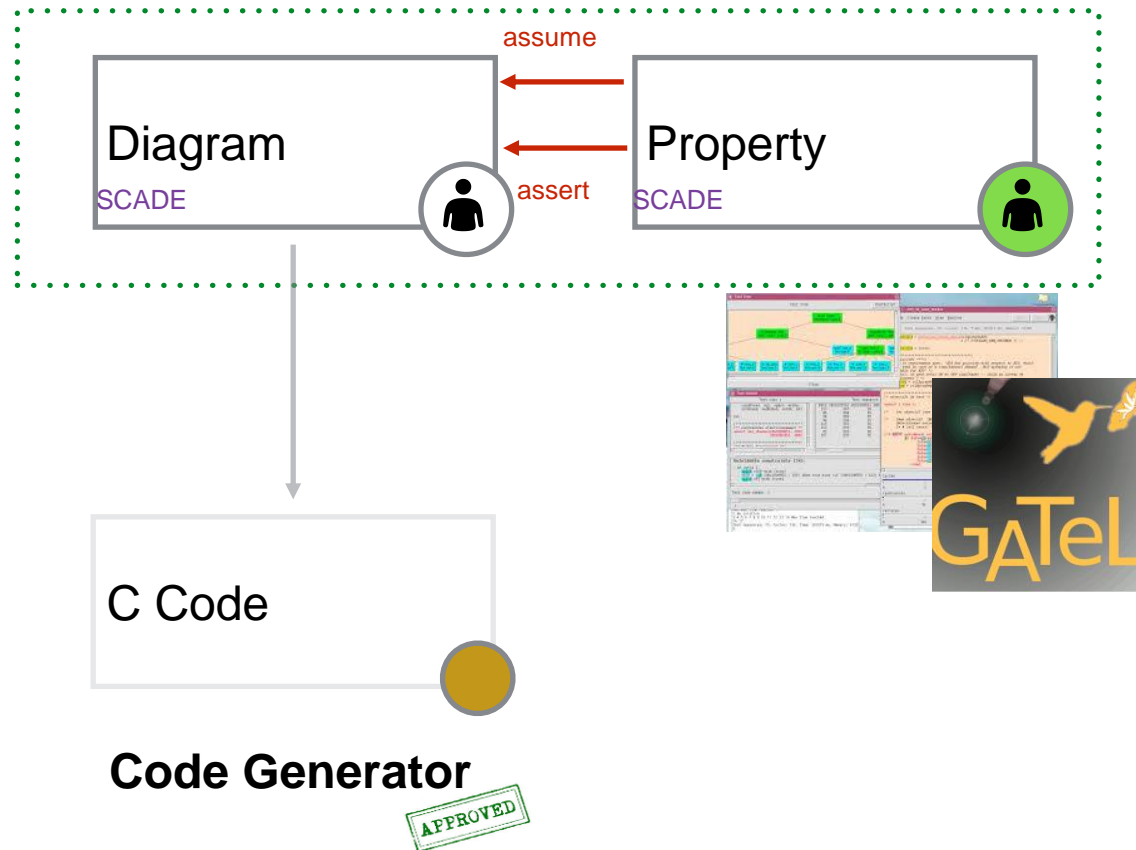



```
int div(int a, int b) {  
  int q;  
  q=0;  
  while (a>=b) {  
    a= a-b;  
    q= q+1;  
  }  
  return q;  
}
```



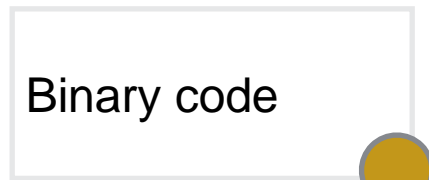
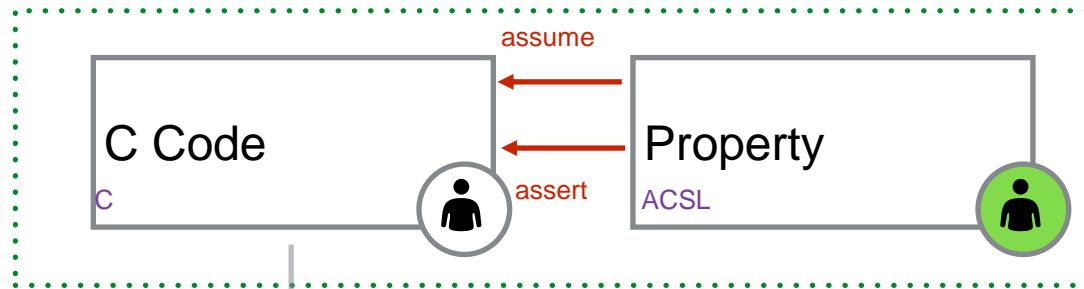
}

VERIFICATION | CEA LIST TOOLS

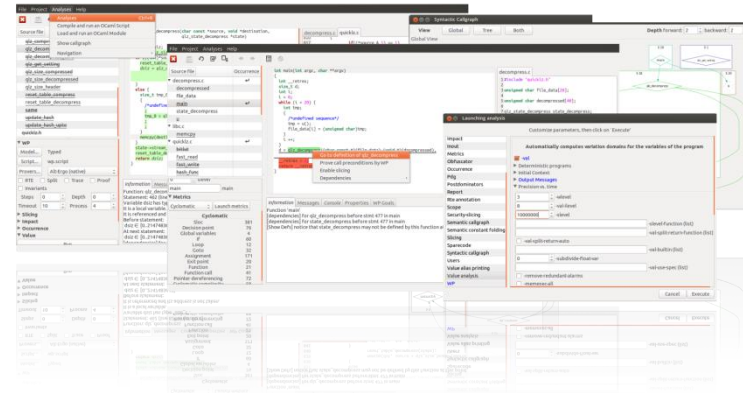




-  Development
-  Generation
-  Validation

VERIFICATION | CEA LIST TOOLS



Compiler



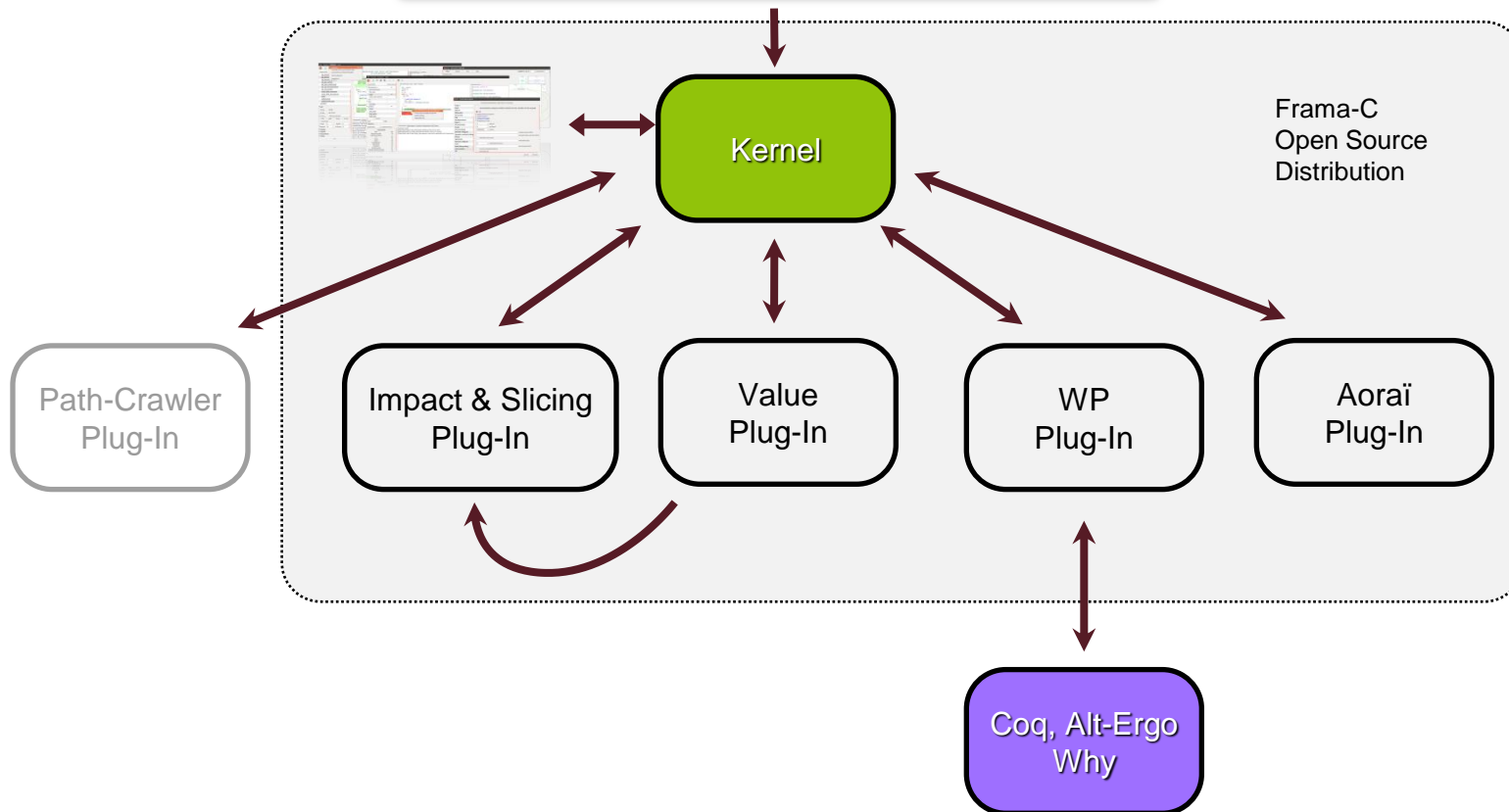
-  Development
-  Generation
-  Validation

VERIFICATION | CORE FRAMA-C COMPONENTS

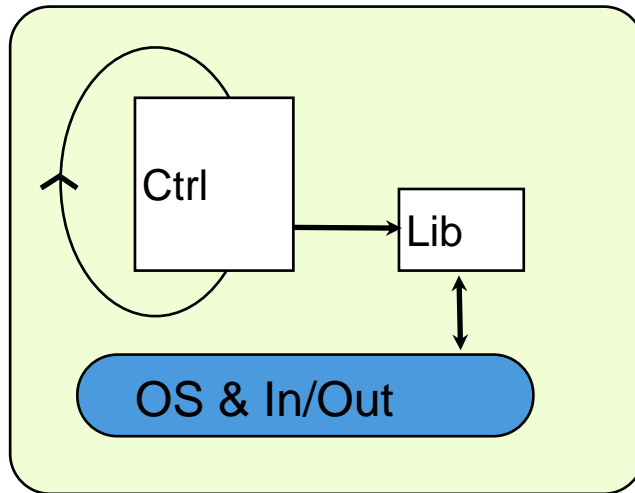
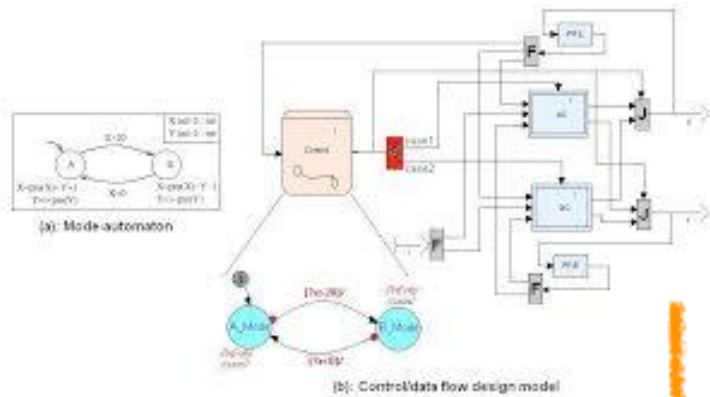
```

/*@ ensures \result >= x && \result >= y;
   ensures \result == x || \result == y;
*/
int max (int x, int y) { return (x > y) ? x : y; }

```



VERIFICATION | STATE OF PRACTICE



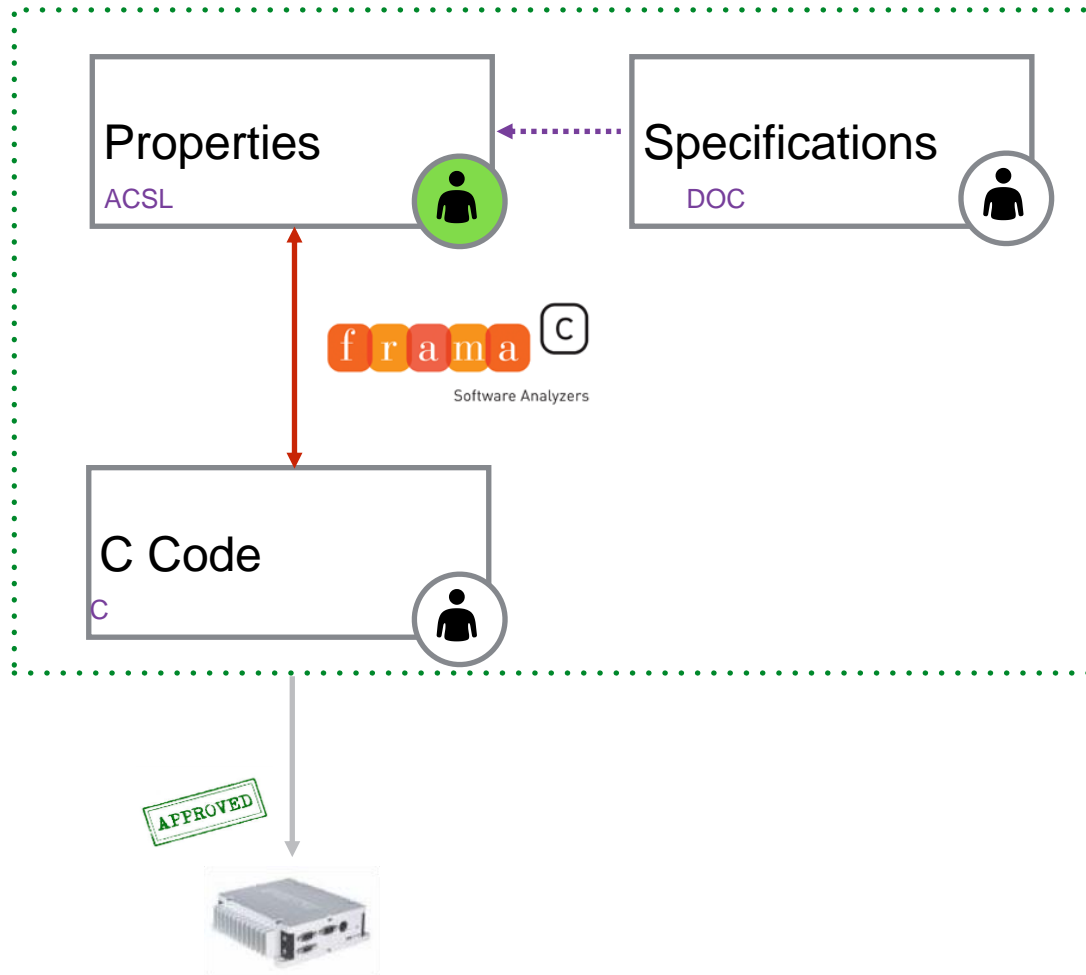
+250 klocs
~10' analysis runtime
~200 remaining alarms



- Value
- WP

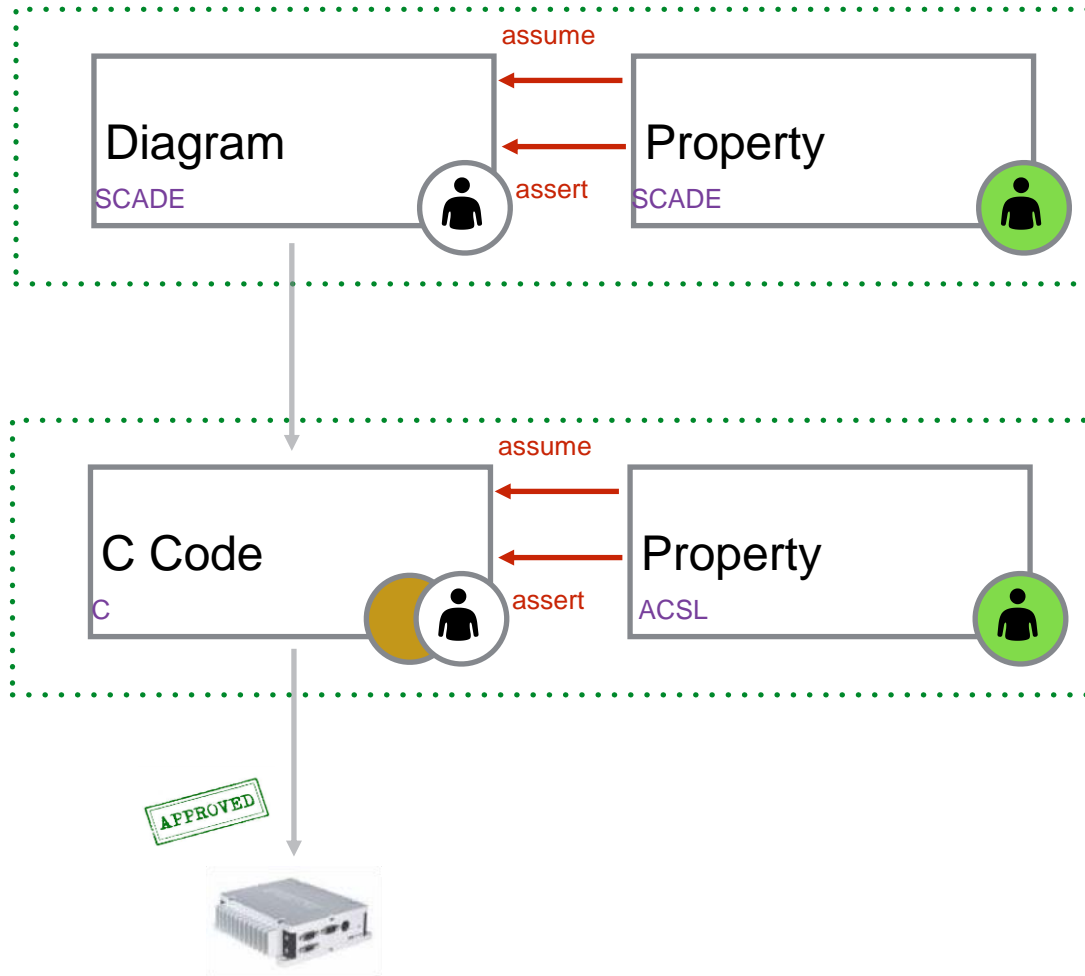
Absence of RTE
OS/Ctrl+Lib separation
Local proofs

VERIFICATION | STATE OF PRACTICE – LOCAL PROOFS



-  Development
-  Generation
-  Validation

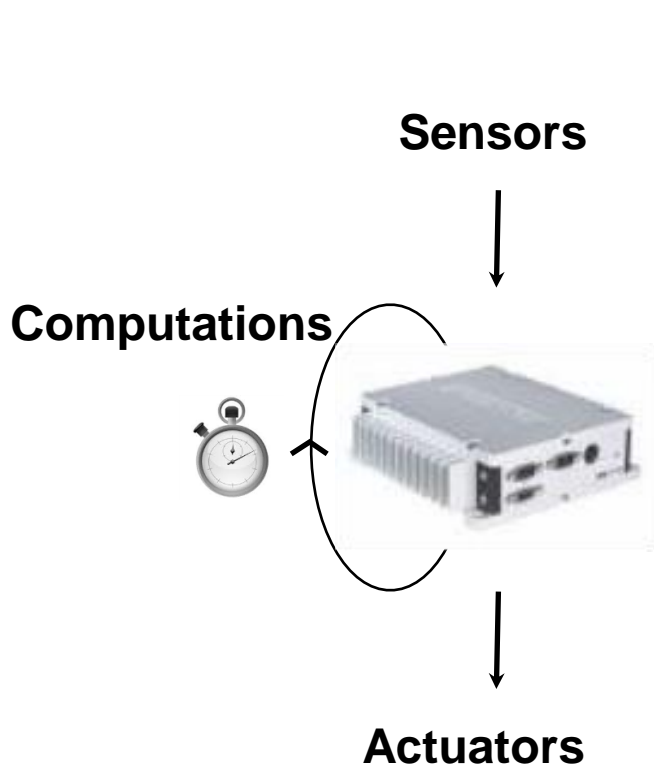
VERIFICATION | MULTIPLE PROOF TECHNIQUES



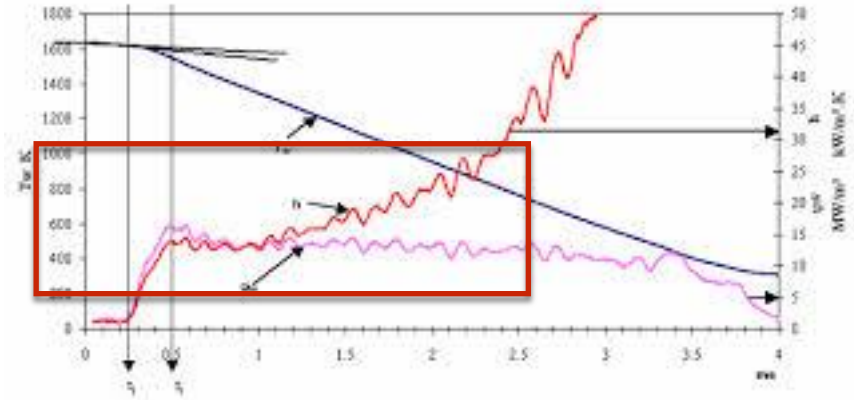
-  Development
-  Generation
-  Validation

}

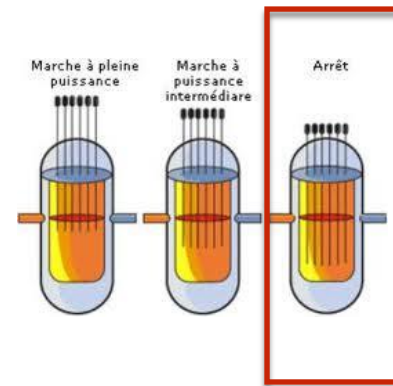
A SPECIAL KIND OF PROPERTY



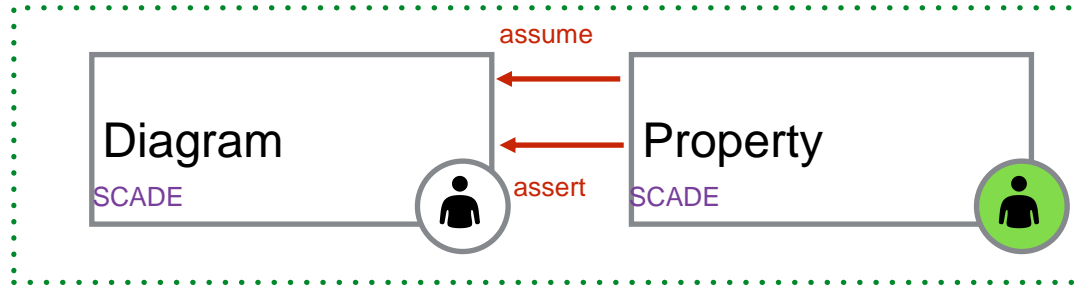
ASSUME



ASSERT



VERIFICATION | DIAGRAM-LEVEL PROPERTIES



$$\bar{y}_t = \varphi(\bar{m}_t, \bar{x}_t)$$

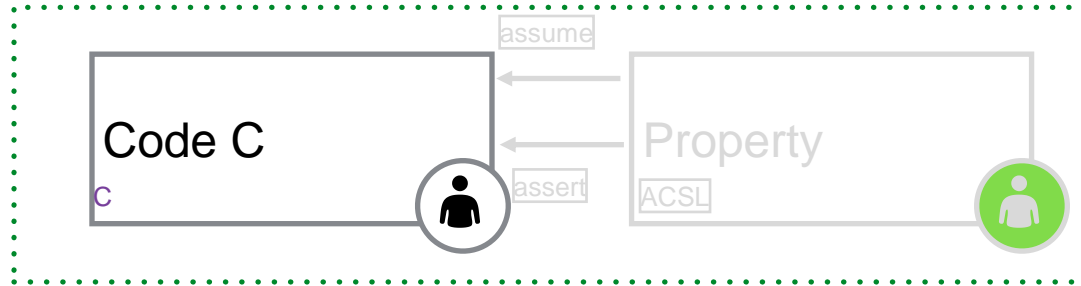
$$\bar{m}_{(t+1)} = M(\bar{m}_t, \bar{x}_t) \quad \bar{a}_{(t+1)} = A(\bar{a}_t, \bar{x}_t, \bar{y}_t)$$

$$\forall (t \leq T), P(\bar{a}_t)$$

$(\varphi, M) = 1$ diagram (abstract level)



VERIFICATION | CODE-LEVEL PROPERTIES

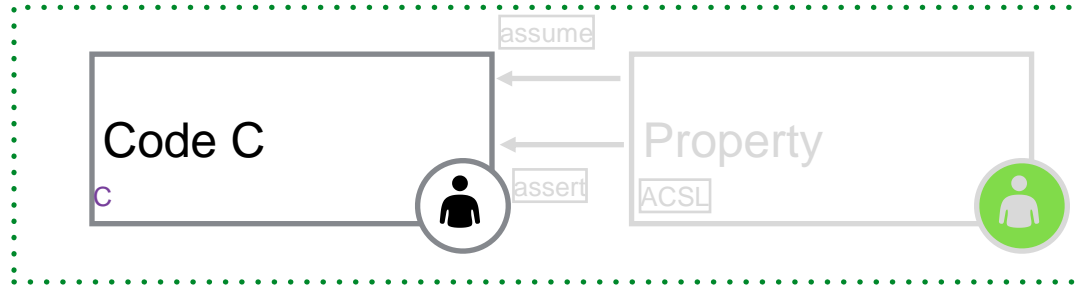


$$\bar{y}_t = \varphi(\bar{m}_t, \bar{x}_t)$$

$$\bar{m}_{(t+1)} = M(\bar{m}_t, \bar{x}_t)$$

$(\varphi, M) = 1$ C function (low level details)

VERIFICATION | CODE-LEVEL PROPERTIES



$$\bar{y}_t = \varphi(\bar{m}_t, \bar{x}_t)$$

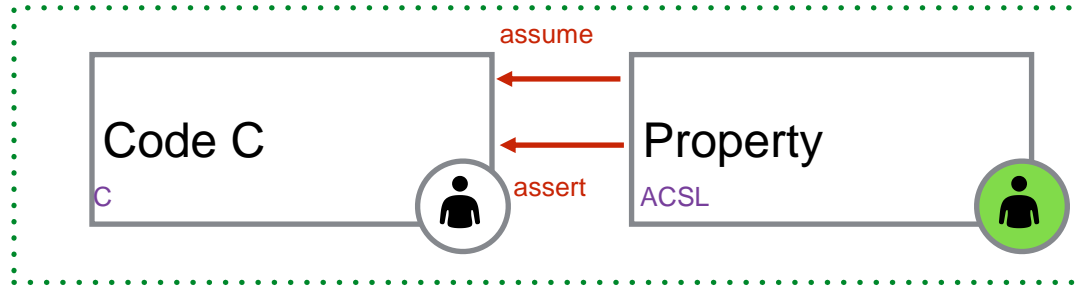
$$\bar{m}_{(t+1)} = M(\bar{m}_t, \bar{x}_t)$$

$$\forall t, (m^i)_t \in D^i$$

(φ, M) = main function (low level details)



VERIFICATION | CODE-LEVEL PROPERTIES



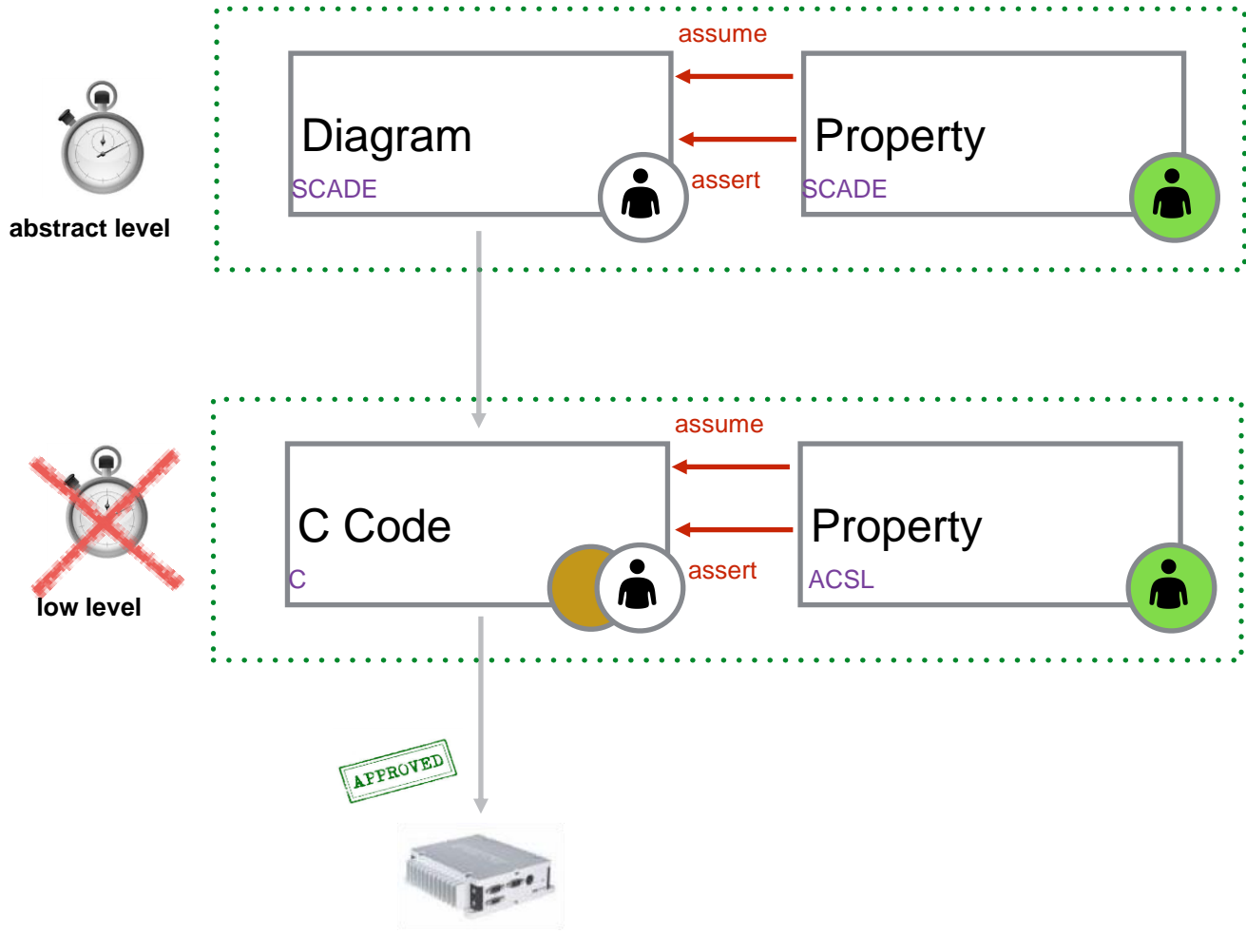
$$\begin{array}{ll} \bar{y}_t = \varphi(\bar{m}_t, \bar{x}_t) & H(\bar{m}_t, \bar{x}_t) \\ \bar{m}_{(t+1)} = M(\bar{m}_t, \bar{x}_t) & G(\bar{m}_{(t+1)}, \bar{x}_t, \bar{y}_t) \end{array}$$

$$\forall t, \quad H(\bar{m}_t, \bar{x}_t) \implies G(\bar{m}_{(t+1)}, \bar{x}_t, \bar{y}_t)$$

$(\varphi, M) = 1$ C function (low level details)

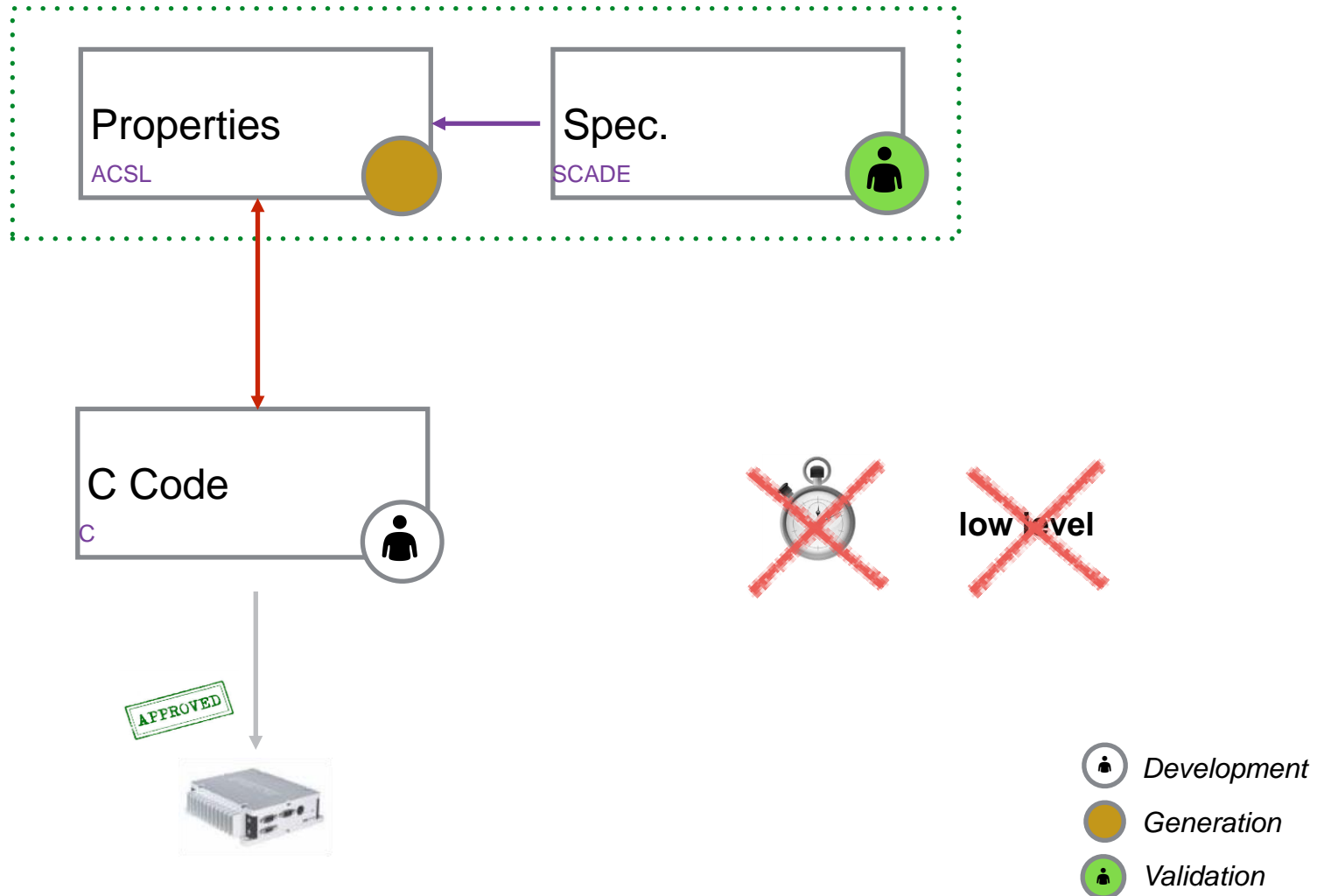


VERIFICATION | DO MORE / DO BETTER?

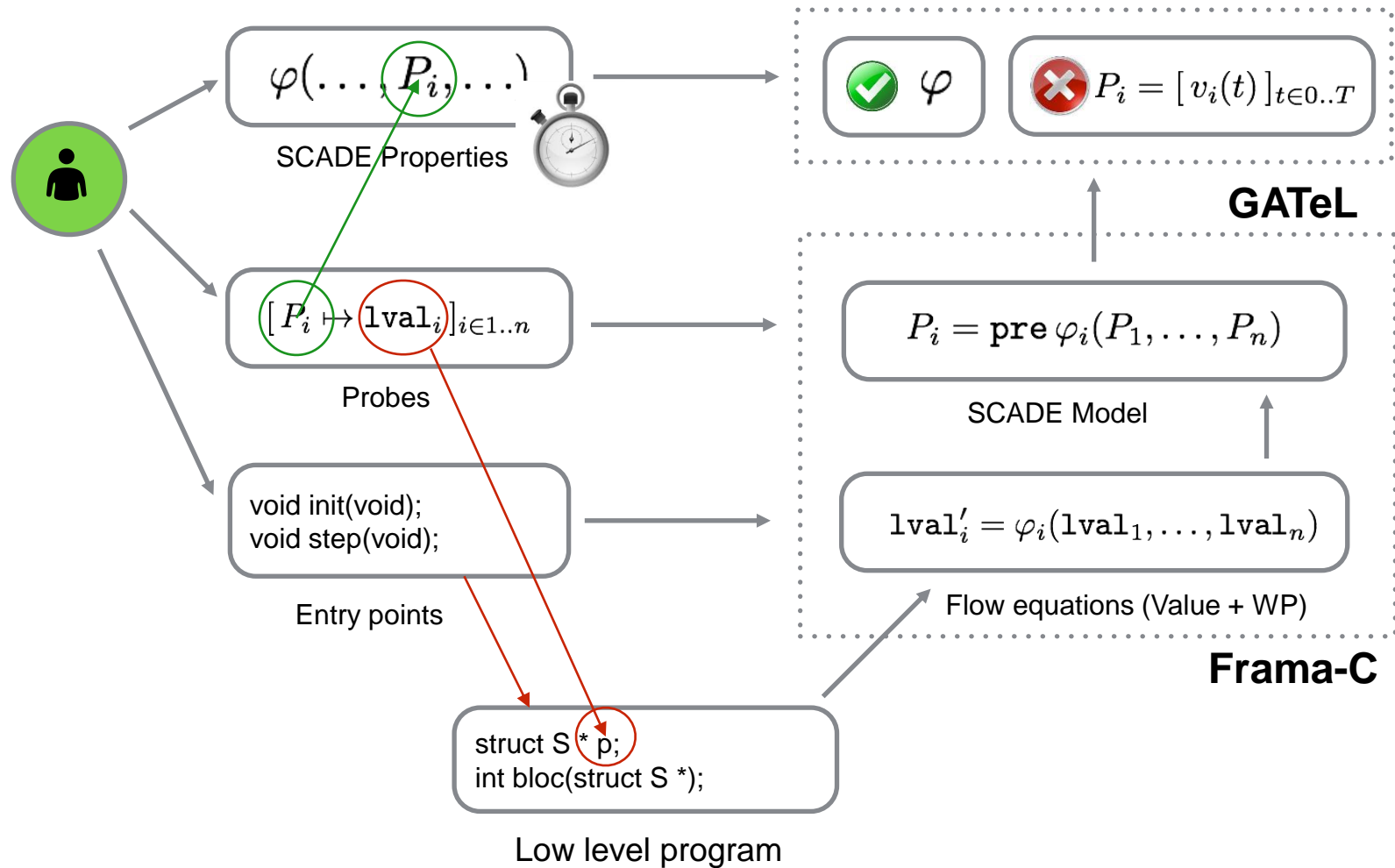


- Development
- Generation
- Validation

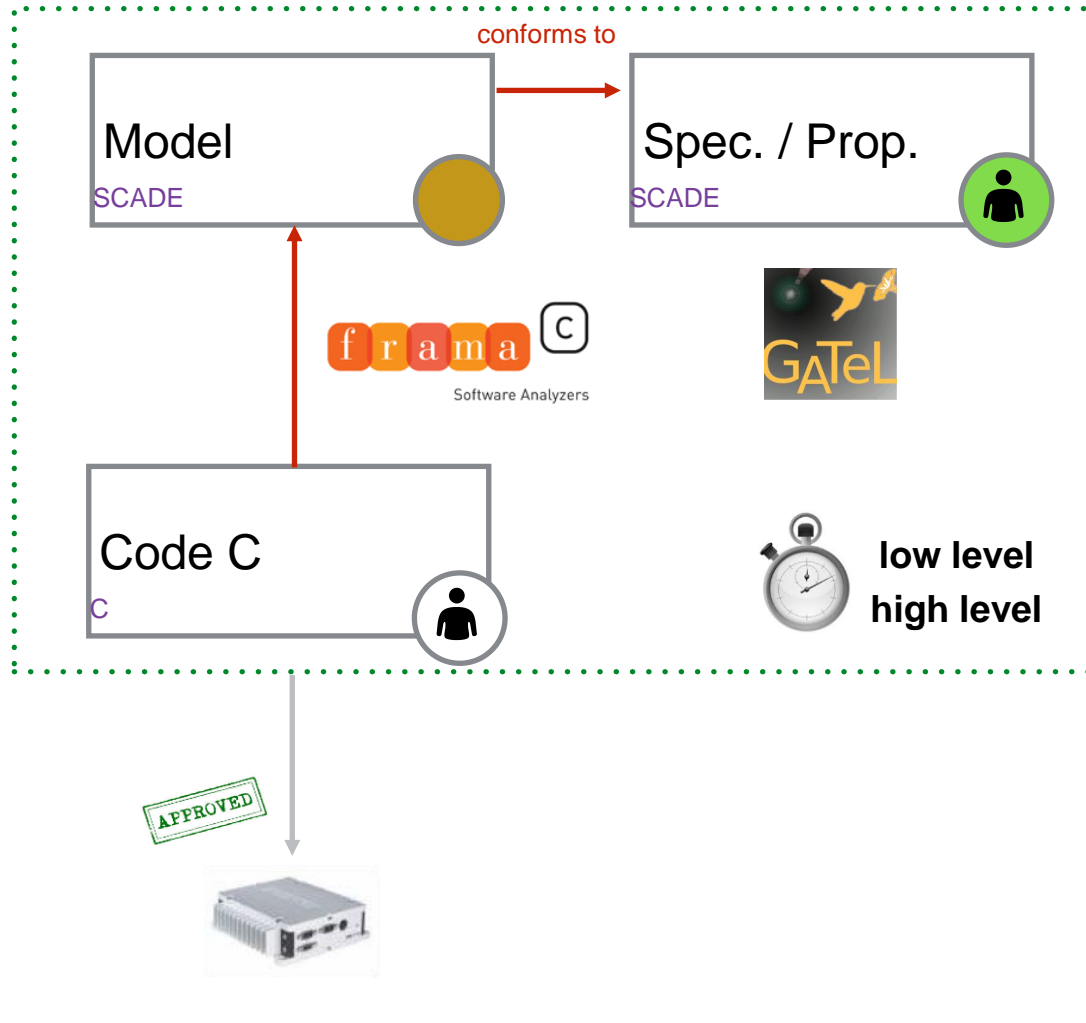
VERIFICATION | PROOFS FROM GENERATED PROPERTIES



VERIFICATION | ANOTHER APPROACH

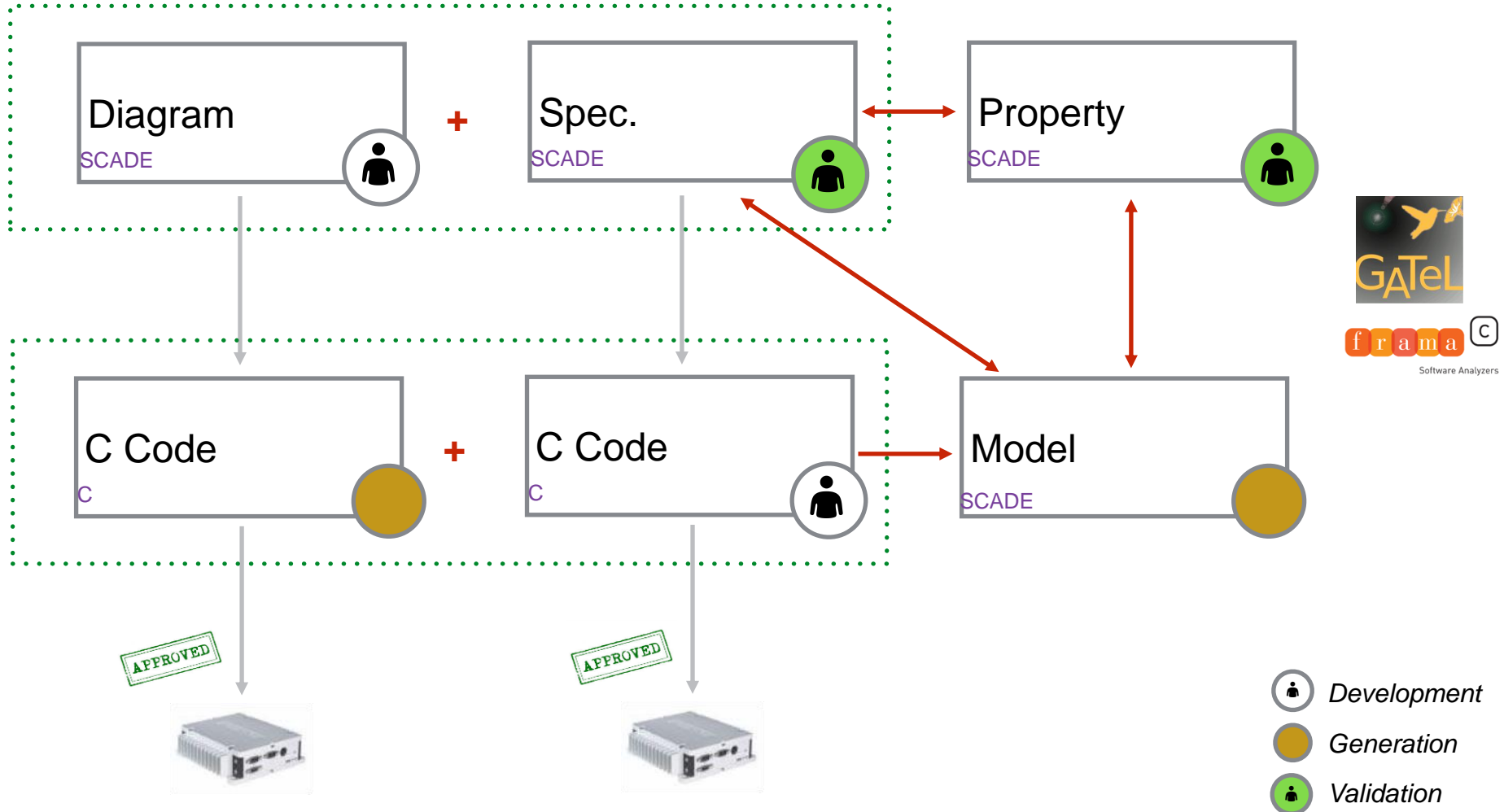


VERIFICATION | ANOTHER APPROACH – 2014

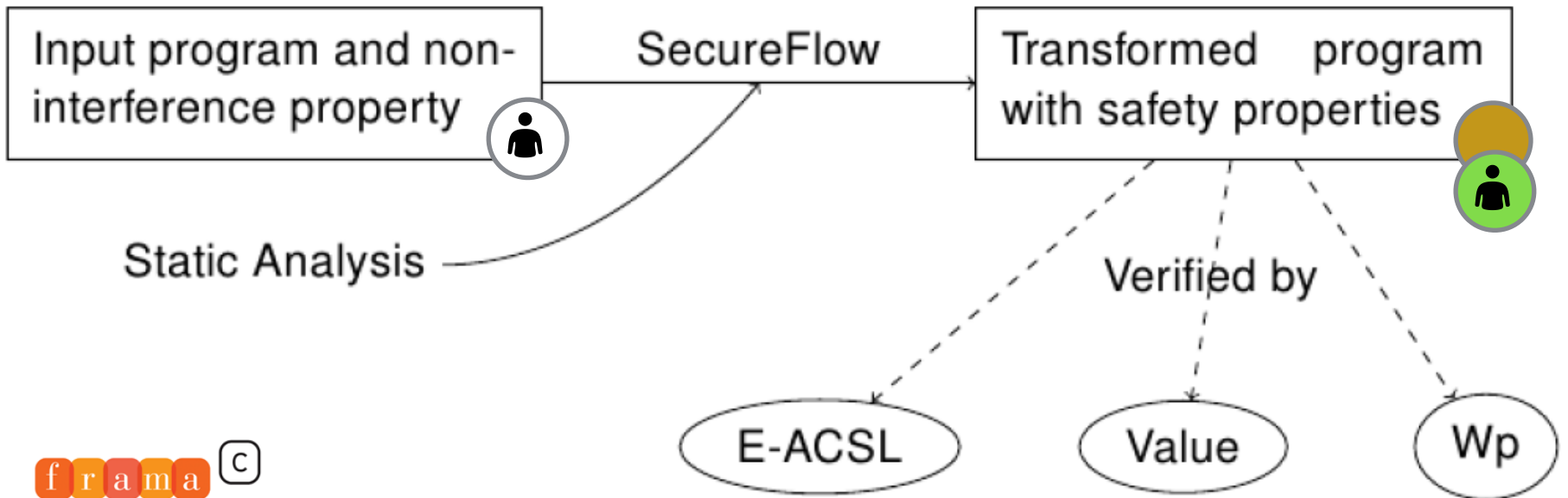


-  Development
-  Generation
-  Validation

CONCLUSION | COMPLEMENTARY NEW APPROACHES

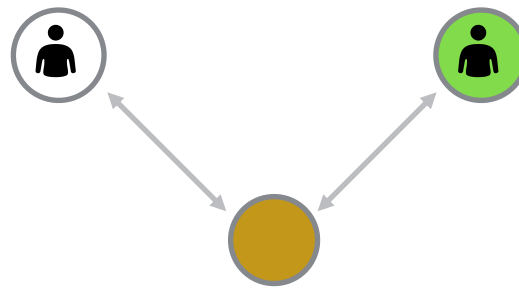


CONCLUSION | A DETOUR THROUGH CYBERSECURITY



CONCLUSION | A NEW DISCIPLINE

Systems engineering ||| Proof engineering



Software Security Laboratory
Software & Systems Engineering Department
CEA LIST

Florent Kirchner
florent.kirchner@cea.fr

This document is the property of CEA. It can not be copied or disseminated without its authorization.