• page 13, line 6: s_2 should read s_3 (twice).

• page 31, equality (2.3.28) in Theorem 2.3.27: in the sum, j should range from 0 to $\min\{k, i\}$, not from $\max\{0, k+i-n\}$ to $\min\{k, i\}$. The reason is that in the expression for $a(k, i, j)$ the second binomial coefficient must be interpreted in the extended sense (cf. page 28, line 3), because in the proof it is shown that the left-hand side and right-hand side are the same as polynomials in the two indeterminates n and x. In the proof on the last line but two "$j < k + i - n$" should be omitted: it is correct since we have assumed that $n > k + i$, but not used, and is misleading.

• page 58, notes of Section 2.6: Theorem 2.6.5 comes from Sloane, Reddy and Chen [596] and

G. V. Zaitsev, V. A. Zinoviev and N. V. Semakov: Interrelation of Preparata and Hamming codes and extension of Hamming codes to new double-error-correcting codes, *Proc. 2nd Internat. Symp. on Information Theory*, Tsahkadsor, September 2-8, 1971. Edited by B. N. Petrov and F. Cski. A complementary volume to Problems of Control and Information Th. / Problemy Upravleniya i Teorii Informatsii. Akadmiai Kiad, Budapest, pp. 257-263, 1973.

• page 65, line -5: "smallest" should read "largest".

• page 74, line 3: "smallest" is redundant.

• page 139, line 10: "in nonzero." should read "is nonzero.".

• page 147, line 4: "Lemma 2.4.6" should read "Theorem 2.4.8".

• page 193, Table 7.1, and page 213, lines -9/-8: $t[12, 6] > 2$ is due to Graham and Sloane [265].

• end of page 229, the linear programming problem: we must add the constraint $\beta_i \geq 0$ for all i.

• page 242, line -7: $m + 2$ should read $m - 2$.

• page 248, Lemma 9.3.6 needs to be restated:

Assume that $m \geq 6$ and $n \geq 2^{m-2}$. Then $R_{RM}(2, m) < t$ if there does not exist a self-complementary $[t, m+1, d \geq t - R_{RM}(1, m-1)]$ code with a generator matrix where all columns are distinct.

• page 250, lines 14/15: if m is odd and the number of columns in $\mathbf{A}$ is $m+1$, there may not be any proper subset $\mathbf{Q}$, but then the sum of all columns is zero and by adding a fixed column to all columns we obtain a zero column which can be deleted.

• page 256, last line: the double inequality should be reversed and read $(h-1)(m-1)/h \leq r \leq m-2$.

• page 258, line -15: > 240 should read $= 240$.

• page 259, line 13: Theorem 9.2.2 should be Theorem 9.2.16.

• page 329, lines 12 and 14:

$$\lim_{n\to\infty} \sup \quad \text{should read} \quad \limsup_{n\to\infty}.$$

• page 353, the value of $R_{10,10}$ is 35, not 34, see J. Carlson and D. Stolarski: The correct solution to Berlekamp's switching game, *Discrete Mathematics*, vol. 287, pp. 145–150, 2004.

• page 408, Lemma 16.3.9: condition (iii) should read
$w(\mathbf{x}) = 2, w(\mathbf{y}) = w(\mathbf{z}) = 3$ and $w(\mathbf{x}+\mathbf{y}) = w(\mathbf{x}+\mathbf{z}) = 1$.

• page 483, about the problem UB-LIN, and pages 486/487, about the construction of an instance of UB-LIN: it should be remarked that the matrix $\mathbf{H}$ in the instance of UB-LIN is of full rank. In the construction, we can assume, without loss of generality, that every element of $X_1 \cup X_2 \cup X_3$ is contained in at least one triple in $M_1 \cup M_2$ (otherwise, the answer is trivially NO). This assumption will guarantee that the constructed matrix $\mathbf{H}$ is of full rank.

• page 504: reference [161] should come before reference [156].