

Coq

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Coq: outline

- Basic notions
 - Introduction
 - History
 - Architecture
 - Vernacular
 - Basic tactics
- Extra features
 - Modules
 - Computation
 - Tactic language
 - Extraction

Outline

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Informations on Coq

- ▶ The Coq web site `coq.inria.fr`
 - ▶ Official distribution + current version of development publically available
 - ▶ Reference manual
 - ▶ Library
 - ▶ User's contributions
- ▶ Reference book : the Coq'art by Yves Bertot and Pierre Castéran



Interactive Theorem Proving and Program Development
Coq'Art: The Calculus of Inductive Constructions
Series: Texts in Theoretical Computer Science.

- ▶ Associated web site with exercises:

<http://www.labri.fr/perso/casteran/CoqArt>

History

See preface of the reference manual

- ▶ 85 (Coquand-Huet) type-checker of the **pure Calculus of Constructions** + **hierarchy of universes**, tactics
- ▶ 88 [V4.10] first version of extraction (Prop/Set) encoded impredicative inductive definitions (Paulin)
- ▶ 89-95 [V5.6] **primitive inductive definitions** with recursor [V5.7] Camllight [V5.8] **(co)-inductive definitions** with pattern-matching and fixpoints (Paulin, Giménez) new architecture (Chet Murthy)
- ▶ 96-01 [V6.1] coercions (A. Saïbi), automatic tactics (Ring, S. Boutin), inversion (C. Cornes) [V6.2] **Efficient reduction** (B. Barras)
- ▶ 02-03 [V7] new functional architecture (Filliâtre), **Modules** (Courant, Chrzęszcz), tactic language (Delahaye), **local definitions** (Herbelin), powerful pattern-matching, new extraction (Letouzey)
- ▶ 04- [V8] **Coq'art**, new syntax, **predicative Set**, **byte-code compilation** (B. Grégoire) ...

Remarkable Developments

- ▶ JavaCard architecture modeling (Trusted Logic, EAL7 certification)
- ▶ Fundamental theorem of Algebra (Constructive proof, CCorn Nijmegen)
- ▶ Four color theorem (Gonthier-Werner)
- ▶ Certified compiler for C (Leroy et al)
- ▶ Primality checker (Théry et al)
- ▶ ...

Architecture

- ▶ Implemented in OCAML
- ▶ Batch compilation : `coqc`
- ▶ Interfaces : `coqide` (`lablgtk`), Proof General (`emacs`), web ...
- ▶ Kernel
 - ▶ strong rules for inductive definitions
complete pattern-matching, guarded definitions
 - ▶ modules
 - ▶ possible efficient byte-code reduction
- ▶ Extensions
 - ▶ universe inference
 - ▶ coercions, implicit arguments
 - ▶ advanced pattern-matching
 - ▶ notations, tactic language, functional definitions
 - ▶ extraction

Basic terms

Prop Set Type

forall $x\ y, B$
 $A \rightarrow B \rightarrow C$

forall $(x\ y : A)\ z, B$
 $(A \rightarrow B) \rightarrow C$

fun $x\ y \Rightarrow t$

fun $(x\ y : A)\ z \Rightarrow t$

$t\ x\ y$

$t\ (u\ x)$

match t **with** .. **end**

let $x := t$ **in** u

fix $f\ (x:A) : B := t$

Basic notations

▶ Propositional logic

$A \wedge B$ $A \vee B$ $A \leftrightarrow B$ $\sim A$

▶ predicate logic

`exists x, P`
`x=y`

▶ data-types

$A * B$ (x, y)

$A + B$

▶ arithmetic (nat, Z ...)

`0` `45` `n+m` `n*m`

Basic Vernacular

Enriching the environments

Definition *name : type := term.*

Variable/Hypothesis *name : type.*

(co)inductive definition

Inductive/CoInductive *name params : type
:= constructors.*

Fixpoint *name params {struct arg}: type := term.*

CoFixpoint *name params : type := term.*

Bottom-up development

Lemma/Theorem *name : type.*

Save/Qed/Defined.

Save/Qed: Opaque constants not unfolded in computation.

Libraries

- ▶ Section mechanism

Section *name*.

Variable ...

Hypothesis ...

Definition ...

Lemma ...

End *name*.

- ▶ Loading precompiled libraries

Require Export *library*.

- ▶ Modules

Basic help

Check *term*.

Print *ident*.

SearchAbout *ident*.

Eval compute **in** *term*.

Print LoadPath.

Basic tactics

- ▶ `exact term`; `assumption`
- ▶ `intros`
- ▶ `apply term`; `constructor`
- ▶ `case term`; `elim term`
- ▶ `change term`; `unfold term`; `simpl`
- ▶ `fix n`; `cofix`

Composed tactics

Inductive types

- ▶ inversion *term*; induction *term*
- ▶ equality : rewrite *term*; injection *term*; discriminate

Automatic tactics

- ▶ Database search: auto, trivial
- ▶ Decision procedure: tauto, firstorder, omega
- ▶ Propositional simplification: intuition
- ▶ Reflexive tactic: ring

Tacticals

- ▶ $tac_1; tac_2$ $tac; [tac_1 | \dots | tac_n]$
- ▶ try $tac, tac_1 || tac_2$

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Modules

- ▶ Extension of Ocaml modules systems to type theory
 - ▶ Modules are not terms
 - ▶ Interfaces contains both abstract and explicit declarations.
 - ▶ Parameterized Modules (functors)
- ▶ Meta-theory by J. Courant (98)
- ▶ Implementation by J. Chrząszcz (02)
- ▶ Non logical features : notations, Hint databases, Extraction are compatible with the module system.

Demo.

Computation

- ▶ Computation is part of type-checking (verification of convertibility)

$$\frac{\Gamma \vdash U : s \quad \Gamma \vdash t : T \quad T \equiv U}{\Gamma \vdash t : U}$$

- ▶ Internal programming language (functional kernel of CAML)
- ▶ Possibility to write complex programs and use them in proofs
 - ▶ Four colors theorem
 - ▶ Reflexive tactics
- ▶ Efficient reduction technics (byte-code compiling, B. Grégoire)
- ▶ Ongoing work on proof-irrelevance

Demo.

Reflexive tactics

- ▶ We have a function $s2bool : \forall AB, A + B \rightarrow bool$ and a proof $strue : \forall p : A + B, s2bool p = true \rightarrow A$
- ▶ If there is a (closed) proof thm of type $\forall x : A, (P x) + Q$ then for closed term a :

$$strue (thm a) (refleq true) : P a$$

Proof is well-typed when $s2bool (thm a) \equiv true$

- ▶ Problems
 - ▶ $a : A$ should be closed (reification), but we also prefer to keep a small
 - ▶ thm should be proved and should reduce efficiently
- ▶ Application
 - ▶ Ring, (R)Omega, Setoid Rewrite...
 - ▶ Interfaces between Coq and other systems using traces.

Example of reflection

```

Inductive form : Set :=
  T | F | Var : nat -> form
  | Conj : form -> form -> form.
Def env := list Prop.
Def find_env (e:env) (n:nat) : Prop := ...
Fixpoint interp (e:env) (f:form) struct f :=
  match f with
    T => True | F => False
  | Conj A B => interp e A /\ interp e B
  | Var n => find_env e n
  end.
Eval compute in
  (interp ((1=1)::(0=0)::nil)
  (Conj (Var 0) (Conj (Var 1) (Var 1))))).
  = 1 = 1 /\ 0 = 0 /\ 0 = 0 : Prop

```

Simplification function

Def `simplform` : `form -> form`.

Lemma `simplcorrect` :

`forall e f, interp e (simplform f) -> interp e f`.

In order to use the simplification

- ▶ Find *e* and *f* such that *interp e f* is convertible with the goal (reification)
- ▶ Apply *simplcorrect*
- ▶ Compute *simplform* and *interp*

Tactic language

Ltac designed by D. Delahaye

- ▶ A way to write complex tactics without writing ML code.
- ▶ A specific language
 - ▶ specific patterns for matching goals (non-linear)

```
match goal with
```

```
  id:?A / ?B |- ?A => case id; trivial  
  | _ => idtac
```

```
end
```

```
match goal with |- context[?a+0] => rewrite ...
```

- ▶ specific notion of backtracking
 - ▶ patterns are tried until the right-hand side succeeds
- ▶ specific constructions : **fresh name**, **type of term** ...

Demo

Extraction

- ▶ Distinction between **Prop** and **Set** given by the user
- ▶ Parts in **Prop** are erased
- ▶ Extraction of Ocaml programs (P. Letouzey)
 - ▶ Termination problems **fun** $x : \perp \Rightarrow t$
 - ▶ Typing problems
 - ▶ polymorphism
 - ▶ generalised inductive types
 - ▶ dependent types : **if** b **then** 0 **else** "hello"
 - ▶ Efficiency

Demo

Coq-lab

Proposed exercises

[http://www.lri.fr/~{}paulin/TypesSummerSchool/
exercises.html](http://www.lri.fr/~{}paulin/TypesSummerSchool/exercises.html)