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TP 6 - Modeling Operational Semantics

Semaine du 8 février 2021

Objective : Defining an operational semantics "Plotkin-style" big-step semantic for as more-or-less standard regular expression language. This type of semantics represents "states" 's in the original language and inductively models the transition relation between states via a transition predicate $\langle _, _ \rangle \longrightarrow_c _$ of type ' $s \times 'output\ option \times 's \Rightarrow bool$ where *option* is the usual option type constructor from `Main`. Following common terminology in automata theory, we will call a list of atoms a *word*, and a set of words a *language*. The *denotational semantics* we are referring here is a function that maps an (abstract) *syntax* to a set of *denotations*, i.e. regular expressions to the language they denote. We will introduce the term *epsilon* as abbreviation for *Star Empty*, an α *rexp* term for the empty language.

We reuse the abstract syntax α *rexp* of the regular expression language from TP5.

Exercice 1 (Inductive Sets - Inductive Proofs)

Define a Plotkin-style semantics for regular expressions, where '*output* is set to ' α . Complete the list of inductive rules starting with :

- $\langle [a], Some\ a \rangle \longrightarrow_c\ \epsilon$
- $\langle Empty : R, None \rangle \longrightarrow_c\ Empty$
- $\langle [a] : R, Some\ a \rangle \longrightarrow_c\ R$
- $\langle Star\ r, None \rangle \longrightarrow_c\ (r : (Star\ r))$
- $\langle Star\ r, None \rangle \longrightarrow_c\ \epsilon$
- ...

Tasks :

1. Prove $\langle [a], Some\ a' \rangle \longrightarrow_c\ \epsilon = (a' = a)$ and $\langle [a], a' \rangle \longrightarrow_c\ R' = (a' = Some\ a \wedge R' = \epsilon)$ (This should hold for your completion of the above inductive rule-set).
2. Derive all similar lemmas resulting from your definitions (should be approx 8). Hint : for the latter rule, there is actually a specific command that derives this type of simplification lemmas automatically. For example, the last mentioned lemma could be derived automatically via :

$$inductive_simps\ atom1S : "\langle [a], a' \rangle \longrightarrow_c\ R"$$

3. Prove the elimination rule :

$$\langle [a], Some\ a' \rangle \longrightarrow_c\ \epsilon \Longrightarrow (a' = a \Longrightarrow P) \Longrightarrow P$$

4. Prove all other elimination rules and configure them into the global context as such. Hint : for the latter rule, there is actually a specific command that derives this type of simplification lemmas automatically. For example, the last mentioned lemma could be derived automatically via :

$$inductive_simps\ atom1S : "\langle [a], a' \rangle \longrightarrow_c\ R"$$

5. Now define the mu ;tiple step semantics Plotkin style. This hould lopok like this :

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inductive
  evalstar :: "[ 'a rexp, 'a list, 'a rexp ] => bool" ("⟨_,_⟩ / →c* _" [0,0,60] 60)
where
  idle:      "⟨epsilon, []⟩ →c* epsilon"
| step1:    "⟨r, Some a⟩ →c r'  =>  ⟨r, [a]⟩ →c* r'"
| continuation1: "⟨r, None⟩ →c r'  =>  ⟨r', S⟩ →c* r'' =>  ⟨r, S⟩ →c* r'''"
| continuation2: "⟨r, Some a⟩ →c r'  =>  ⟨r', S⟩ →c* r'' =>  ⟨r, a#S⟩ →c* r'''"

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6. Prove :

$$\langle Star((\lfloor CHR''a'' \rfloor \parallel \lfloor CHR''b'' \rfloor) : \lfloor CHR''c'' \rfloor), "bc'' \rangle \rightarrow_c^* \epsilon$$

7. Prove :

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theorem operational_implies_denotational_generalized':
assumes nat_steps: "⟨r, s⟩ →c* r'"
and den_cont: "∃ as. as ∈ L r'"
shows "∃ as . s@as ∈ L r ∧ as ∈ L r'"

```

8. Prove the main theorem "operational semantics implies denotational semantics" :

$$\langle r, s \rangle \rightarrow_c^* \epsilon \rightarrow s \in L(r)$$

Note : Main provides the notation CHR ''a'' for "the character a". Strings are defined as lists of characters.

Exercise 2 (OPTIONAL : Report)

(IN CASE THAT YOU WANT TO HAVE IT GRADED. RECALL THAT 2 OUT OF 6 TP's SHOULD BE SUBMITTED.)

1. Write a little report answering all questions above, note the difficulties you met, add some screenshots if appropriate. 3 pages max (except screenshots and other figures).